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Undecidable propositional bimodal logics and one-variable first-order linear temporal logics with counting

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Abstract

First-order temporal logics are notorious for their bad computational behaviour. It is known that even the two-variable monadic fragment is highly undecidable over various linear timelines, and over branching time even one-variable fragments might be undecidable. However, there have been several attempts on finding well-behaved fragments of first-order temporal logics and related temporal description logics, mostly either by restricting the available quantifier patterns, or considering sub-Boolean languages. Here we analyse seemingly ‘mild’ extensions of decidable one-variable fragments with counting capabilities, interpreted in models with constant, decreasing, and expanding first-order domains. We show that over most classes of linear orders these logics are (sometimes highly) undecidable, even without constant and function symbols, and with the sole temporal operator ‘eventually’.

We establish connections with bimodal logics over 2D product structures having linear and ‘difference’ (inequality) component relations, and prove our results in this bimodal setting. We show a general result saying that satisfiability over many classes of bimodal models with commuting ‘unbounded’ linear and difference relations is undecidable. As a by-product, we also obtain new examples of finitely axiomatisable but Kripke incomplete bimodal logics. Our results generalise similar lower bounds on bimodal logics over products of two linear relations, and our proof methods are quite different from the known proofs of these results. Unlike previous proofs that first ‘diagonally encode’ an infinite grid, and then use reductions of tiling or Turing machine problems, here we make direct use of the grid-like structure of product frames and obtain lower complexity bounds by reductions of counter (Minsky) machine problems. Representing counter machine runs apparently requires less control over neighbouring grid-points than tilings or Turing machine runs, and so this technique is possibly more versatile, even if one component of the underlying product structures is ‘close to’ being the universal relation.

1 Introduction

1.1 First-order linear temporal logic with counting.

Though first-order temporal logics are natural and expressive languages for querying and constraining temporal databases [7, 8] and reasoning about knowledge that changes in time [25], their practical use has been discouraged by their high computational complexity. It is well-known that even the two-variable monadic fragment is undecidable over various linear timelines, and its satisfiability problem is Σ_1^1 -hard over the natural numbers [47, 48, 35, 12, 13]. Also, even the one-variable fragment of first-order branching time logic CTL^* is undecidable [26]. Still, similarly to classical first-order logic where the decision problems of its fragments were studied and classified in great detail [5], there have been a number of attempts on finding the border between decidable and undecidable fragments of first-order temporal logics and related temporal description logics, mostly either by restricting the available quantifier patterns [8, 24, 25, 3, 9, 21, 22, 31], or considering sub-Boolean languages [30, 2].

In this paper we contribute to this ‘classificational’ research line by considering seemingly ‘mild’ extensions of decidable one-variable fragments. We study the satisfiability problem of the one-variable ‘future’ fragment of linear temporal logic with counting to two, interpreted in models over various timelines, and having constant, decreasing, or expanding first-order domains. Our language FOLTL^\neq keeps all Boolean connectives, it has no restriction on formula-generation, and it is strong enough to express uniqueness of a property of domain elements ($\exists^=1x$), and the ‘elsewhere’ quantifier ($\forall^\neq x$). However, FOLTL^\neq -formulas use only a single variable (and so contain only monadic predicate symbols), FOLTL^\neq has no equality, no constant or function symbols, and its only temporal operators are ‘eventually’ and ‘always in the future’. FOLTL^\neq is weaker than the two-variable monadic *monodic* fragment with equality, where temporal operators can be applied only to subformulas with at most one free variable. (This fragment with the ‘next time’ operator is known to be Σ_1^1 -hard over the natural numbers [50, 10].) FOLTL^\neq is connected to bimodal product logics [14, 13] (see also below), and to the temporalisation of the expressive description logic \mathcal{CQ} with one global universal role [49]. Here are some examples of FOLTL^\neq -formulas:

- “An order can only be submitted once.” $\forall x \Box_F (\text{Subm}(x) \rightarrow \Box_F \neg \text{Subm}(x))$.
- The Barcan formula: $\exists x \Diamond_F P(x) \leftrightarrow \Diamond_F \exists x P(x)$.
- “Every day has its unique dog.” $\Box_F \exists^=1 x \text{Dog}(x) \wedge \Box_F \forall x (\text{Dog}(x) \rightarrow \Box_F \neg \text{Dog}(x))$.
- “It’s only me who is always unlucky.” $\Box_F \neg \text{Lucky}(x) \wedge \forall^\neq x \Diamond_F \text{Lucky}(x)$.

Note that FOLTL^\neq can also be considered as a fragment of three-variable classical first-order logic with only binary predicate symbols, but it is not within the guarded fragment.

Our contribution While the addition of ‘elsewhere’ quantifiers to the two-variable fragment of classical first-order logic does not increase the NEXPTIME complexity of its satisfiability problem [17, 18, 37], we show that adding the same feature to the (decidable) one-variable fragment of first-order temporal logic results in (sometimes highly) undecidable logics over most linear timelines, not only in models with constant domains, but even those with decreasing and expanding first-order domains. Our main results on the FOLTL^\neq -satisfiability problem are summarised in Fig. 1.

	$\langle \omega, < \rangle$	all finite linear orders	all linear orders	$\langle \mathbb{Q}, < \rangle$ or $\langle \mathbb{R}, < \rangle$
constant domains	Σ_1^1 -complete Cor. 3.4	undecidable r.e. Cor. 3.4	undecidable co-r.e. Cor. 4.3	undecidable Cor. 4.17
decreasing domains	Σ_1^1 -complete Cor. 3.4	undecidable r.e. Cor. 3.4	undecidable co-r.e. Cor. 4.18	undecidable Cor. 4.17
expanding domains	undecidable co-r.e. Cors. 5.2, 5.15	Ackermann-hard decidable Cors. 5.4, 5.17	<div style="border: 1px solid black; padding: 2px;">decidable?</div> co-r.e. Cor. 5.13	<div style="border: 1px solid black; padding: 2px;">decidable?</div>

Figure 1: FOLTL^\neq -satisfiability over various timelines and first-order domains.

1.2 Bimodal logics and two-dimensional modal logics.

It is well-known that the first-order quantifier $\forall x$ can be considered as an ‘**S5**-box’: a propositional modal necessity operator interpreted over relational structures $\langle W, R \rangle$ where $R = W \times W$ (*universal frames*, in modal logic parlance). Therefore, the two-variable fragment of classical first-order logic is related to propositional bimodal logic over two-dimensional (2D) *product frames* [33]. Similarly, the ‘elsewhere’ quantifier $\forall^{\neq} x$ can be regarded as a ‘**Diff**-box’: a propositional modal necessity operator interpreted over *difference frames* $\langle W, \neq \rangle$ where \neq is the inequality relation on W . Looking at FOLTL $^{\neq}$ this way, it turns out that it is just a notational variant of the propositional bimodal logic over 2D products of linear orders and difference frames (Prop. 2.3).

Propositional multimodal languages interpreted in various product-like structures show up in many other contexts, and connected to several other multi-dimensional logical formalisms, such as modal and temporal description logics, and spatio-temporal logics (see [13, 28] for surveys and references). The product construction as a general combination method on modal logics was introduced in [43, 45, 14], and has been extensively studied ever since.

Our contribution We study the *satisfiability problem* of our logics in the propositional bimodal setting. We show that satisfiability over many classes of bimodal frames with commuting linear and difference relations are undecidable (Theorems 3.2, 4.1), sometimes not even recursively enumerable (Theorems 3.1, 4.11). As a by-product, we also obtain new examples of finitely axiomatisable but *Kripke incomplete* bimodal logics (Cor. 4.13). It is easy to see (Prop. 2.2) that satisfiability over *decreasing* or *expanding* subframes of product frames is always reducible to ‘full rectangular’ product frame-satisfiability. We show cases when expanding frame-satisfiability is genuinely simpler than product-satisfiability (Theorems 5.14, 5.16), while it is still very complex (Theorems 5.3, 5.1).

Our findings are in sharp contrast with the much lower complexity of bimodal logics over products of linear and universal frames: Satisfiability over these is usually decidable with complexity between EXPSpace and 2EXPTIME [23, 38]. In particular, we answer negatively a question of [38] by showing that the addition of the ‘horizontal’ difference operator to the decidable 2D product of Priorian Temporal Logic over the class of all linear orders and **S5** results in an undecidable logic (Cor. 4.2).

Our lower bound results are also interesting because they seem to be proper generalisations of similar results about modal products where both components are linear [32, 39, 16, 15, 27]. Satisfiability over linear and difference frames is of the same (NP-complete) complexity, and so there are reductions from ‘linear-satisfiability’ to ‘difference-satisfiability’ and vice versa. However, while we show (Section 5.2) how to ‘lift’ some ‘difference to linear’ reduction to the 2D level, one cannot hope for such a lifting of a reverse ‘linear to difference’ reduction: Satisfiability over ‘difference \times difference’ type products is decidable (being a fragment of two-variable classical first-order logic with counting), while ‘linear \times linear’-satisfiability is undecidable [39].

Our undecidability proofs are quite different from most known undecidability proofs about 2D product logics with transitive components [32, 39, 15]. Even if frames with two commuting relations (and so product frames) always have grid-like substructures, there are two issues one needs to deal with in order to encode grid-based complex problems into them:

- to generate infinity, and
- somehow to ‘access’ or ‘refer to’ neighbouring-grid points, even when there might be further non-grid points around, there is no ‘next-time’ operator in the language, and the relations are transitive and/or dense and/or even ‘close to’ universal.

Unlike previous proofs that first ‘diagonally encode’ the $\omega \times \omega$ -grid, and then use reductions of tiling or Turing machine problems, here we make direct use of the grid-like substructures in commutative frames, and obtain lower bounds by reductions of counter (Minsky) machine problems. Representing counter machine runs apparently requires less control over neighbouring grid-points than tilings or Turing machine runs, and so this technique is possibly more versatile (see Section 2.5 for more details).

Structure Section 2 provides all the necessary definitions, and establishes connections between the two different formalisms. All results are then proved in the propositional bimodal setting. In particular, Section 3 deals with the constant and decreasing domain cases over $\langle \omega, < \rangle$ and finite linear orders. More general results on bimodal logics with ‘linear’ and ‘difference’ components are in Section 4. The expanding domain cases are treated in Section 5. Finally, in Section 6 we discuss some related open problems.

Some of the results appeared in the extended abstract [19].

2 Preliminaries

2.1 Propositional bimodal logics

Below we introduce all the necessary notions and notation. For more information on bimodal logics, consult e.g. [4, 13].

We define *bimodal formulas* by the following grammar:

$$\phi ::= P \mid \neg\phi \mid \phi \wedge \psi \mid \Diamond_0\phi \mid \Diamond_1\phi$$

where P ranges over an infinite set of propositional variables. We use the usual abbreviations \vee , \rightarrow , \leftrightarrow , $\perp := P \wedge \neg P$, $\top := \neg\perp$, $\Box_i := \neg\Diamond_i\neg$, and also

$$\Diamond_i^+\phi := \phi \vee \Diamond_i\phi, \quad \Box_i^+\phi := \phi \wedge \Box_i\phi,$$

for $i = 0, 1$. For any bimodal formula ϕ , we denote by $\text{sub}\phi$ the set of its subformulas.

A *2-frame* is a tuple $\mathfrak{F} = \langle W, R_0, R_1 \rangle$ where R_i are binary relations on the non-empty set W . A *model based on \mathfrak{F}* is a pair $\mathfrak{M} = (\mathfrak{F}, \nu)$, where ν is a function mapping propositional variables to subsets of W . The *truth relation* $\mathfrak{M}, w \models \phi$ is defined, for all $w \in W$, by induction on ϕ as follows:

- $\mathfrak{M}, w \models P$ iff $w \in \nu(P)$,
- $\mathfrak{M}, w \models \neg\phi$ iff $\mathfrak{M}, w \not\models \phi$, $\mathfrak{M}, w \models \phi \wedge \psi$ iff $\mathfrak{M}, w \models \phi$ and $\mathfrak{M}, w \models \psi$,
- $\mathfrak{M}, w \models \Diamond_i\phi$ iff there exists $v \in W$ such that wR_iv and $\mathfrak{M}, v \models \phi$ (for $i = 0, 1$).

We say that ϕ is *satisfied in \mathfrak{M}* , if there is $w \in W$ with $\mathfrak{M}, w \models \phi$. Given a set Σ of bimodal formulas, we write $\mathfrak{M} \models \Sigma$ if we have $\mathfrak{M}, w \models \phi$, for every $\phi \in \Sigma$ and every $w \in W$. We say that ϕ is *valid in \mathfrak{F}* , if $\mathfrak{M}, w \models \phi$, for every model \mathfrak{M} based on \mathfrak{F} and for every $w \in W$. If every formula in a set Σ is valid in \mathfrak{F} , then we say that \mathfrak{F} is a *frame for Σ* . We let $\text{Fr}\Sigma$ denote the class of all frames for Σ .

A set L of bimodal formulas is called a (normal) *bimodal logic* (or *logic*, for short) if it contains all propositional tautologies and the formulas $\Box_i(p \rightarrow q) \rightarrow (\Box_ip \rightarrow \Box_iq)$, for $i = 0, 1$, and is closed under the rules of Substitution, Modus Ponens and Necessitation $\varphi/\Box_i\varphi$, for $i = 0, 1$. Given a bimodal logic L , we will consider the following problem:

L-SATISFIABILITY: Given a bimodal formula ϕ , is there a model \mathfrak{M} such that $\mathfrak{M} \models L$ and ϕ is satisfied in \mathfrak{M} ?

For any class \mathcal{C} of 2-frames, we always obtain a logic by taking

$$\text{Log}\mathcal{C} = \{\phi : \phi \text{ is a bimodal formula valid in every member of } \mathcal{C}\}.$$

We say that $\text{Log}\mathcal{C}$ is *determined by \mathcal{C}* , and call such a logic *Kripke complete*. (We write just $\text{Log}\mathfrak{F}$ for $\text{Log}\{\mathfrak{F}\}$.) Clearly, if $L = \text{Log}\mathcal{C}$, then there might exist frames for L that are not in \mathcal{C} , but L -satisfiability is the same as the following problem:

C-SATISFIABILITY: Given a bimodal formula ϕ , is there a 2-frame $\mathfrak{F} \in \mathcal{C}$ such that ϕ is satisfied in a model based on \mathfrak{F} ?

Commutators and products We might regard bimodal logics as ‘combinations’ of their unimodal¹ ‘components’. Let L_0 and L_1 be two unimodal logics formulated using the same propositional variables and Booleans, but having different modal operators (\Diamond_0 for L_0 and \Diamond_1 for L_1). Their *fusion* $L_0 \oplus L_1$ is the smallest bimodal logic that contains both L_0 and L_1 . The *commutator* $[L_0, L_1]$ of L_0 and L_1 is the smallest bimodal logic that contains $L_0 \oplus L_1$ and the formulas

$$\Box_1 \Box_0 P \rightarrow \Box_0 \Box_1 P, \quad \Box_0 \Box_1 P \rightarrow \Box_1 \Box_0 P, \quad \Diamond_0 \Box_1 P \rightarrow \Box_1 \Diamond_0 P. \quad (1)$$

Commutators are introduced in [14], where it is also shown that a 2-frame $\langle W, R_0, R_1 \rangle$ validates the formulas (1) iff

- R_0 and R_1 *commute*: $\forall x, y, z (xR_0 y R_1 z \rightarrow \exists u (xR_1 u R_0 z))$, and
- R_0 and R_1 are *confluent*: $\forall x, y, z (xR_0 y \wedge xR_1 z \rightarrow \exists u (yR_1 u \wedge zR_0 u))$.

Note that if at least one of R_0 or R_1 is symmetric, then confluence follows from commutativity.

Next, we introduce some special ‘two-dimensional’ 2-frames for commutators. Given frames $\mathfrak{F}_0 = \langle W_0, R_0 \rangle$ and $\mathfrak{F}_1 = \langle W_1, R_1 \rangle$, their *product* is defined to be the 2-frame

$$\mathfrak{F}_0 \times \mathfrak{F}_1 = \langle W_0 \times W_1, \tilde{R}_0, \tilde{R}_1 \rangle,$$

where $W_0 \times W_1$ is the Cartesian product of W_0 and W_1 and, for all $u, u' \in W_0, v, v' \in W_1$,

$$\begin{aligned} \langle u, v \rangle \tilde{R}_0 \langle u', v' \rangle & \text{ iff } uR_0 u' \text{ and } v = v', \\ \langle u, v \rangle \tilde{R}_1 \langle u', v' \rangle & \text{ iff } vR_1 v' \text{ and } u = u'. \end{aligned}$$

2-frames of this form will be called *product frames* throughout. For classes \mathcal{C}_0 and \mathcal{C}_1 of unimodal frames, we define

$$\mathcal{C}_0 \times \mathcal{C}_1 = \{\mathfrak{F}_0 \times \mathfrak{F}_1 : \mathfrak{F}_i \in \mathcal{C}_i, \text{ for } i = 0, 1\}.$$

Now, for $i = 0, 1$, let L_i be a Kripke complete unimodal logic in the language with \Diamond_i . The *product* of L_0 and L_1 is defined as the (Kripke complete) bimodal logic

$$L_0 \times L_1 = \text{Log}(\text{Fr } L_0 \times \text{Fr } L_1).$$

Product frames always validate the formulas in (1), and so it is not hard to see that $[L_0, L_1] \subseteq L_0 \times L_1$ always holds. If both L_0 and L_1 are Horn axiomatisable, then $[L_0, L_1] = L_0 \times L_1$ [14]. In general, $[L_0, L_1]$ can not only be properly contained in $L_0 \times L_1$, but there might even be infinitely many logics in between [29, 20].

The following result of Gabbay and Shehtman [14] is one of the few general ‘transfer’ results on the satisfiability problem of 2D logics. It is an easy consequence of the recursive enumerability of the consequence relation of classical (many-sorted) first-order logic:

Theorem 2.1. *If \mathcal{C}_0 and \mathcal{C}_1 are classes of frames such that both are recursively first-order definable in the language having a binary predicate symbol, then $\mathcal{C}_0 \times \mathcal{C}_1$ -satisfiability is co-r.e., that is, its complement is recursively enumerable.*

Expanding and decreasing 2-frames Product frames are special cases of the following construction for getting 2D frames. Take a (‘horizontal’) frame $\mathfrak{F} = \langle W, R \rangle$ and a sequence $\overline{\mathfrak{G}} = \langle \mathfrak{G}_u = \langle W_u, R_u \rangle : u \in W \rangle$ of (‘vertical’) frames. We can define a 2-frame by taking

$$\mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}} = \langle \{ \langle u, v \rangle : u \in W, v \in W_u \}, \tilde{R}_0, \tilde{R}_1 \rangle,$$

where

$$\begin{aligned} \langle u, v \rangle \tilde{R}_0 \langle u', v' \rangle & \text{ iff } uRu' \text{ and } v = v', \\ \langle u, v \rangle \tilde{R}_1 \langle u', v' \rangle & \text{ iff } vR_u v' \text{ and } u = u'. \end{aligned}$$

Clearly, if $\mathfrak{G}_x = \mathfrak{G}_y = \mathfrak{G}$ for all x, y in \mathfrak{F} , then $\mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}} = \mathfrak{F} \times \mathfrak{G}$. However, we can put slightly milder assumptions on the \mathfrak{G}_x . We call a 2-frame of the form $\mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}}$

¹Syntax and semantics of *unimodal* logics are defined similarly to bimodal ones, using only one of the two modal operators. Throughout, 1-frames will be called simply *frames*.

- an *expanding 2-frame* if \mathfrak{G}_x is a subframe² of \mathfrak{G}_y whenever xRy , and
- a *decreasing 2-frame* if \mathfrak{G}_y is a subframe of \mathfrak{G}_x whenever xRy .

So product frames are both expanding and decreasing 2-frames. Expanding 2-frames always validate $\Box_0\Box_1P \rightarrow \Box_1\Box_0P$ and $\Diamond_0\Box_1P \rightarrow \Box_1\Diamond_0P$ (but not necessarily $\Box_1\Box_0P \rightarrow \Box_0\Box_1P$), and decreasing 2-frames validate $\Box_1\Box_0P \rightarrow \Box_0\Box_1P$ (but not necessarily the other two formulas in (1)).

For classes \mathcal{C}_0 and \mathcal{C}_1 of frames, we define

$$\begin{aligned}\mathcal{C}_0 \times^e \mathcal{C}_1 &= \{\text{expanding 2-frame } \mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}} : \mathfrak{F} \in \mathcal{C}_0, \mathfrak{G}_x \in \mathcal{C}_1 \text{ for all } x \text{ in } \mathfrak{F}\}, \\ \mathcal{C}_0 \times^d \mathcal{C}_1 &= \{\text{decreasing 2-frame } \mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}} : \mathfrak{F} \in \mathcal{C}_0, \mathfrak{G}_x \in \mathcal{C}_1 \text{ for all } x \text{ in } \mathfrak{F}\}.\end{aligned}$$

It is not hard to see that for all classes $\mathcal{C}_0, \mathcal{C}_1$ of frames, both $\mathcal{C}_0 \times^d \mathcal{C}_1$ -satisfiability and $\mathcal{C}_0 \times^e \mathcal{C}_1$ -satisfiability is reducible to $\mathcal{C}_0 \times \mathcal{C}_1$ -satisfiability. Indeed, take a fresh propositional variable D (for *domain*), and for every bimodal formula ϕ , define ϕ^D by relativising each occurrence of \Diamond_0 and \Diamond_1 in ϕ to D . Let n be the nesting depth of the modal operators in ϕ , any for any formula ψ and $i = 0, 1$, let

$$\Box_i^{\leq n} \psi := \bigwedge_{k \leq n} \overbrace{\Box_i \dots \Box_i}^k \psi.$$

Then we have (cf. [13, Thm.9.12]):

Proposition 2.2.

- ϕ is $\mathcal{C}_0 \times^d \mathcal{C}_1$ -satisfiable iff $D \wedge \Box_0^{\leq n} \Box_1^{\leq n} (\Diamond_0 D \rightarrow D) \wedge \phi^D$ is $\mathcal{C}_0 \times \mathcal{C}_1$ -satisfiable.
- ϕ is $\mathcal{C}_0 \times^e \mathcal{C}_1$ -satisfiable iff $D \wedge \Box_0^{\leq n} \Box_1^{\leq n} (D \rightarrow \Box_0 D) \wedge \phi^D$ is $\mathcal{C}_0 \times \mathcal{C}_1$ -satisfiable.

‘Linear’ and ‘difference’ logics Throughout, a frame $\langle W, R \rangle$ is called *rooted* with root $r \in W$ if every $w \in W$ can be reached from r by taking finitely many R -steps. By a *linear order* we mean an irreflexive³, transitive and trichotomous relation. Let \mathcal{C}_{lin} and \mathcal{C}_{lin}^{fin} denote the classes of all linear orders and all finite linear orders, respectively. We let $\mathbf{K4.3} := \text{Log } \mathcal{C}_{lin}$, that is, the unimodal logic determined by all linear orders. $\mathbf{K4.3}$ is well-studied as a temporal logic, and it is well-known that frames for $\mathbf{K4.3}$ are *weak orders*.⁴ A linear order $\langle W, R \rangle$ is called a *well-order* if every non-empty subset of W has an R -least element.

We denote by \mathcal{C}_{diff} (\mathcal{C}_{diff}^{fin}) the class of all (finite) *difference frames*, that is, frames of the form $\langle W, \neq \rangle$ where \neq is the inequality relation on W . We let $\mathbf{Diff} := \text{Log } \mathcal{C}_{diff}$, that is, the unimodal logic determined by all difference frames. From the axiomatisation of \mathbf{Diff} by Segerberg [44] it follows that frames for \mathbf{Diff} are *pseudo-equivalence*⁵ relations. If \mathfrak{M} is a model based on a rooted pseudo-equivalence frame, then we can express the uniqueness of a modally definable property in \mathfrak{M} . For any formula ϕ ,

$$\Diamond^=1 \phi := \Diamond^+(\phi \wedge \Box \neg \phi).$$

Then, $\Diamond^=1 \phi$ is satisfied in \mathfrak{M} iff there is a unique w with $\mathfrak{M}, w \models \phi$.

As all the axioms of $\mathbf{K4.3}$ and \mathbf{Diff} , and the formulas in (1) are Sahlqvist formulas, the commutator $[\mathbf{K4.3}, \mathbf{Diff}]$ is Sahlqvist axiomatisable, and so Kripke complete. Also,

$$\begin{aligned}\text{Fr}[\mathbf{K4.3}, \mathbf{Diff}] &= \{\langle W, R_0, R_1 \rangle : R_0 \text{ is a weak order,} \\ &\quad R_1 \text{ is a pseudo-equivalence, } R_0 \text{ and } R_1 \text{ commute}\} \quad (2)\end{aligned}$$

(for more information on Sahlqvist formulas and canonicity, consult e.g. [4, 6]).

² $\langle W, R \rangle$ is called a *subframe* of $\langle U, S \rangle$, if $W \subseteq U$ and $R = S \cap (W \times W)$.

³This is just for simplifying the overall presentation. Reflexive cases are covered in Section 4.3.

⁴A relation R is called a *weak order* if it is transitive and *weakly connected*: $\forall x, y, z (xRy \wedge xRz \rightarrow (y = z \vee yRz \vee zRy))$. In other words, a rooted weak order is a linear chain of clusters of universally connected points.

⁵A relation R is called a *pseudo-equivalence* if it is symmetric and *pseudo-transitive*: $\forall x, y, z (xRyRz \rightarrow (x = z \vee xRz))$. So a pseudo-equivalence is almost an equivalence relation, just it might have both reflexive and irreflexive points.

2.2 One-variable first-order linear temporal logic with counting to two

We define FOLTL[≠]-formulas by the following grammar:

$$\phi ::= P(x) \mid \neg\phi \mid \phi \wedge \psi \mid \Diamond_F \phi \mid \exists^{\neq} x \phi$$

where (with a slight abuse of notation) P ranges over an infinite set \mathcal{P} of *monadic* predicate symbols.

A FOLTL-model is a tuple $\mathfrak{M} = \langle \langle T, < \rangle, D_t, I \rangle_{t \in T}$, where $\langle T, < \rangle$ is a linear order, representing the *timeline*, D_t is a non-empty set, the *domain at moment* t , for each $t \in T$, and I is a function associating with every $t \in T$ a first-order structure $I(t) = \langle D_t, P^{I(t)} \rangle_{P \in \mathcal{P}}$. We say that \mathfrak{M} is *based on* the linear order $\langle T, < \rangle$. \mathfrak{M} is a *constant* (resp. *decreasing*, *expanding*) *domain model*, if $D_t = D_{t'}$, (resp. $D_t \supseteq D_{t'}$, $D_t \subseteq D_{t'}$) whenever $t, t' \in T$ and $t < t'$. A constant domain model is clearly both a decreasing and expanding domain model as well, and can be represented as a triple $\langle \langle T, < \rangle, D, I \rangle$.

The *truth-relation* $(\mathfrak{M}, t) \models^a \phi$ (or simply $t \models^a \phi$ if \mathfrak{M} is understood) is defined, for all $t \in T$ and $a \in D_t$, by induction on ϕ as follows:

- $t \models^a P(x)$ iff $a \in P^{I(t)}$, $t \models^a \neg\phi$ iff $t \not\models^a \phi$, $t \models^a \phi \wedge \psi$ iff $t \models^a \phi$ and $t \models^a \psi$,
- $t \models^a \exists^{\neq} x \phi$ iff there exists $b \in D_t$ such that $b \neq a$ and $t \models^b \phi$,
- $t \models^a \Diamond_F \phi$ iff there is $t' \in T$ such that $t' > t$, $a \in D_{t'}$ and $t' \models^a \phi$.

We say that ϕ is *satisfiable in* \mathfrak{M} if $\mathfrak{M}, t \models^a \phi$ holds for some $t \in T$ and $a \in D_t$. Given a class \mathcal{C} of linear orders, we say that ϕ is FOLTL[≠]-satisfiable in *constant* (*decreasing*, *expanding*) *domain models over* \mathcal{C} , if ϕ is satisfiable in some constant (*decreasing*, *expanding*) domain FOLTL-model based on some linear order from \mathcal{C} .

We introduce the following abbreviations:

$$\exists x \phi := \phi \vee \exists^{\neq} x \phi, \quad \exists^{\geq 2} x \phi := \exists x (\phi \wedge \exists^{\neq} x \phi).$$

It is straightforward to see that they have the intended semantics:

- $t \models^a \exists x \phi$ iff there exists $b \in D_t$ with $t \models^b \phi$,
- $t \models^a \exists^{\geq 2} x \phi$ iff there exist $b, b' \in D_t$ with $b \neq b'$, $t \models^b \phi$ and $t \models^{b'} \phi$.

Also, we could have chosen $\exists x$ and $\exists^{\geq 2} x$ as our primary connectives instead of $\exists^{\neq} x$, as

$$\exists^{\neq} x \phi \leftrightarrow (\neg\phi \wedge \exists x \phi) \vee \exists^{\geq 2} x \phi.$$

2.3 Connections between propositional bimodal logic and FOLTL[≠]

Clearly, one can define a bijection $*$ from FOLTL[≠]-formulas to bimodal formulas, mapping each $P(x)$ to P , $\Diamond_F \phi$ to $\Diamond_0 \phi^*$, $\exists^{\neq} x \phi$ to $\Diamond_1 \phi^*$, and commuting with the Booleans. Also, there is a bijection † between constant domain FOLTL-models $\mathfrak{M} = \langle \langle T, < \rangle, D, I \rangle$ and modal models $\mathfrak{M}^\dagger = \langle \mathfrak{F}, \nu \rangle$ where $\mathfrak{F} = \langle T, < \rangle \times \langle D, \neq \rangle$ and $\nu(P) = \{ \langle t, a \rangle : \mathfrak{M}, t \models^a P(x) \}$. Similarly, there is a one-to-one connection between expanding (*decreasing*) 2-frames with linear ‘horizontal’ and difference ‘vertical’ components, and expanding (*decreasing*) domain FOLTL-models. So it is straightforward to see the following:

Proposition 2.3. *For any class \mathcal{C} of linear orders, and any FOLTL[≠]-formula ϕ ,*

- ϕ is FOLTL[≠]-satisfiable in constant domain models over \mathcal{C} iff ϕ^* is $\mathcal{C} \times \mathcal{C}_{\text{diff}}$ -satisfiable;
- ϕ is FOLTL[≠]-satisfiable in expanding domain models over \mathcal{C} iff ϕ^* is $\mathcal{C} \times {}^e \mathcal{C}_{\text{diff}}$ -satisfiable;
- ϕ is FOLTL[≠]-satisfiable in decreasing domain models over \mathcal{C} iff ϕ^* is $\mathcal{C} \times {}^d \mathcal{C}_{\text{diff}}$ -satisfiable.

2.4 Counter machines

A *Minsky* or *counter machine* M is described by a finite set Q of states, a set $H \subseteq Q$ of terminal states, a finite set $C = \{c_0, \dots, c_{N-1}\}$ of counters with $N > 1$, a finite nonempty set $I_q \subseteq Op_C \times Q$ of instructions, for each $q \in Q - H$, where each operation in Op_C is one of the following forms, for some $i < N$:

- c_i^{++} (increment counter c_i by one),
- c_i^{--} (decrement counter c_i by one),
- $c_i^{??}$ (test whether counter c_i is zero).

A *configuration* of M is a tuple $\langle q, \mathbf{c} \rangle$ with $q \in Q$ representing the current state, and an N -tuple $\mathbf{c} = \langle c_0, \dots, c_{N-1} \rangle$ of natural numbers representing the current contents of the counters. For each $\iota \in Op_C$, we say that there is a (*reliable*) ι -step between configurations $\sigma = \langle q, \mathbf{c} \rangle$ and $\sigma' = \langle q', \mathbf{c}' \rangle$ (written $\sigma \rightarrow^\iota \sigma'$) iff there is $\langle \iota, q' \rangle \in I_q$ such that

- either $\iota = c_i^{++}$ and $c'_i = c_i + 1$, $c'_j = c_j$ for $j \neq i$, $j < N$,
- or $\iota = c_i^{--}$ and $c_i > 0$, $c'_i = c_i - 1$, $c'_j = c_j$ for $j \neq i$, $j < N$,
- or $\iota = c_i^{??}$ and $c'_i = c_i = 0$, $c'_j = c_j$ for $j < N$.

We write $\sigma \rightarrow \sigma'$ iff $\sigma \rightarrow^\iota \sigma'$ for some $\iota \in Op_C$. For each $\iota \in Op_C$, we write $\sigma \rightarrow_{lossy}^\iota \sigma'$ if there are configurations $\sigma^1 = \langle q, \mathbf{c}^1 \rangle$ and $\sigma^2 = \langle q', \mathbf{c}^2 \rangle$ such that $\sigma^1 \rightarrow^\iota \sigma^2$, $c_i \geq c_i^1$ and $c_i^2 \geq c'_i$ for every $i < N$. We write $\sigma \rightarrow_{lossy} \sigma'$ iff $\sigma \rightarrow_{lossy}^\iota \sigma'$ for some $\iota \in Op_C$. A sequence $\langle \sigma_n : n < B \rangle$ of configurations, with $0 < B \leq \omega$, is called a *run* (resp. *lossy run*), if $\sigma_{n-1} \rightarrow \sigma_n$ (resp. $\sigma_{n-1} \rightarrow_{lossy} \sigma_n$) holds for every $0 < n < B$.

Below we list the counter machine problems we will use in our lower bound proofs.

CM NON-TERMINATION: (Π_1^0 -hard [36])

Given a counter machine M and a state q_0 , does M have an infinite run starting with $\langle q_0, \mathbf{0} \rangle$?

CM REACHABILITY: (Σ_1^0 -hard [36])

Given a counter machine M , a configuration $\sigma_0 = \langle q_0, \mathbf{0} \rangle$ and a state q_r , does M have a run starting with σ_0 and reaching q_r ?

CM RECURRENCE: (Σ_1^1 -hard [1])

Given a counter machine M and two states q_0, q_r , does M have a run starting with $\langle q_0, \mathbf{0} \rangle$ and visiting q_r infinitely often?

LCM REACHABILITY: (Ackermann-hard [41])

Given a counter machine M , a configuration $\sigma_0 = \langle q_0, \mathbf{0} \rangle$ and a state q_r , does M have a lossy run starting with σ_0 and reaching q_r ?

The Ackermann-hardness of this problem is shown by Schnoebelen [41] without the restriction that σ_0 has all-0 counters. It is not hard to see that this restriction does not matter: For every M and σ_0 one can define a machine M^{σ_0} that first performs incrementation steps filling the counters up to their ' σ_0 -level', and then performs M 's actions. Then M has a lossy run starting with σ_0 and reaching q_r iff M^{σ_0} has a lossy run starting with all-0 counters and reaching q_r .

LCM ω -REACHABILITY: (Π_1^0 -hard [27, 34, 40])

Given a counter machine M , a configuration $\sigma_0 = \langle q_0, \mathbf{0} \rangle$ and a state q_r , is it the case that for every $n < \omega$ M has a lossy run starting with σ_0 and visiting q_r at least n times?

2.5 Representing counter machine runs in our logics

Before stating and proving our results, here we give a short informal guide on how we intend to use counter machines in the various lower bound proofs of the paper. To begin with, using two different propositional variables S (for *state*) and N (for *next*), we force a ‘diagonal staircase’ with the following properties:

- (i) every S -point ‘vertically’ (R_1) sees some N -point, and
- (ii) every N -point has an S -point as its ‘immediate horizontal (R_0) successor’.

This way we not only force infinity, but also get a ‘horizontal’ next-time operator:

$$X\phi := \Box_1(N \rightarrow \Box_0(S \rightarrow \phi))$$

(see Fig. 2). In the simplest case of product frames of the form $\langle \omega, < \rangle \times \langle W, \neq \rangle$, a grid-like structure with subsequent columns comes by definition, so everything is ready for encoding counter machine runs in them: Subsequent states of a run will be represented by subsequently generated S -points, and the content of each counter c_i at step n of a run will be represented by the number of C_i -points at the n th column of the grid, for some formula C_i (see Fig. 2). As in difference frames uniqueness of a property is modally expressible, we can faithfully express the subsequent changes of the counters (see Section 3).

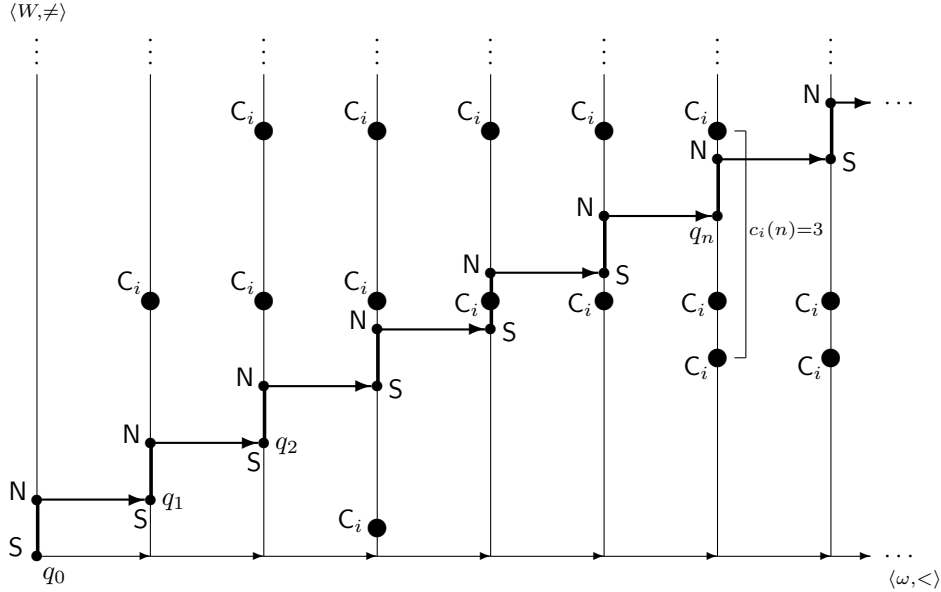


Figure 2: Representing counter machine runs in product frames $\langle \omega, < \rangle \times \langle D, \neq \rangle$ ‘going forward’.

When generalising this technique to ‘timelines’ other than $\langle \omega, < \rangle$, there can be additional difficulties. Say, (ii) above is clearly not doable over dense linear orders. Instead of working with R_0 -connected points, we work with ‘ R_0 -intervals’ and have the ‘interval-analogue’ of (ii): Every N -interval has an S -interval as its ‘immediate R_0 -successor’ (see Section 4.3).

We also generalise our results not only to decreasing 2-frames but for more ‘abstract’ 2-frames having commuting weak order and pseudo-equivance relations (see (2)). In the abstract case, we face an additional difficulty: While commutativity does force the presence of grid-points once a diagonal staircase is present, there might be many other non-grid points in the corresponding ‘vertical columns’, so the control over runs becomes more complicated. In these cases, both the diagonal staircase and counter machine runs are forced going ‘backward’ (see Fig. 3), as this way seemingly gives us greater control over the ‘intended’ grid-points (see Section 4.1).

CLAIM 3.5. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid}$. Then there exists an infinite sequence $\langle y_m \in W : m < \omega \rangle$ of points such that, for all $m < \omega$,*

- (i) $y_0 = r$ and for all $n < m$, $y_m \neq y_n$,
- (ii) $\mathfrak{M}, \langle m, y_m \rangle \models S$,
- (iii) if $m > 0$ then $\mathfrak{M}, \langle m-1, y_m \rangle \models N$.

Proof. By induction on m . To begin with, $\mathfrak{M}, \langle 0, y_0 \rangle \models S$ by (3). Now suppose that for some $m < \omega$ we have $\langle y_k : k \leq m \rangle$ as required. As $\mathfrak{M}, \langle m, y_m \rangle \models S$ by the IH, by (4) there is y_{m+1} such that $\mathfrak{M}, \langle m, y_{m+1} \rangle \models N$. We have $y_{m+1} \neq y_n$ for $n \leq m$ by (3), (5) and the IH. Finally, $\mathfrak{M}, \langle m+1, y_{m+1} \rangle \models S$ follows by (5). \square

Given a counter machine M , we will encode runs that start with all-0 counters by going forward along the created diagonal staircase. For each counter $i < N$, we take two fresh propositional variables C_i^+ and C_i^- . At each moment n of time, these will be used to mark those pairs $\langle n, \dots \rangle$ in \mathfrak{M} where M increments and decrements counter c_i at step n . The actual content of counter c_i is represented by those pairs $\langle n, \dots \rangle$ where $C_i^+ \wedge \neg C_i^-$ holds. The following formula ensures that each ‘vertical coordinate’ in \mathfrak{M} is used only once, and only previously incremented points can be decremented:

$$\text{counter} := \bigwedge_{i < N} \square_0^+ \square_1^+ ((C_i^+ \rightarrow \square_0 C_i^+) \wedge (C_i^- \rightarrow \square_0 C_i^-) \wedge (C_i^- \rightarrow C_i^+)).$$

For each $i < N$, the following formulas simulate the possible changes in the counters:

$$\begin{aligned} \text{Fix}_i &:= \square_1^+ (\square_0 C_i^+ \rightarrow C_i^+) \wedge \square_1^+ (\square_0 C_i^- \rightarrow C_i^-), \\ \text{Inc}_i &:= \diamond_1^+ (\neg C_i^+ \wedge \square_0 C_i^+) \wedge \square_1^+ (\square_0 C_i^- \rightarrow C_i^-), \\ \text{Dec}_i &:= \diamond_1^+ (\neg C_i^- \wedge \square_0 C_i^-) \wedge \square_1 (\square_0 C_i^+ \rightarrow C_i^+). \end{aligned}$$

It is straightforward to prove the following:

CLAIM 3.6. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid} \wedge \text{counter}$ and let, for all $m < \omega$, $i < N$, $c_i(m) := |\{w \in W : \mathfrak{M}, \langle m, w \rangle \models C_i^+ \wedge \neg C_i^-\}|$. Then*

$$c_i(m+1) = \begin{cases} c_i(m), & \text{if } \mathfrak{M}, \langle m, y_m \rangle \models \text{Fix}_i, \\ c_i(m) + 1, & \text{if } \mathfrak{M}, \langle m, y_m \rangle \models \text{Inc}_i, \\ c_i(m) - 1, & \text{if } \mathfrak{M}, \langle m, y_m \rangle \models \text{Dec}_i. \end{cases}$$

Using the above machinery, we can encode the various counter machine instructions. For each $\iota \in \text{Op}_C$, we define the formula Do_ι by taking

$$\text{Do}_\iota := \begin{cases} \text{Inc}_i \wedge \bigwedge_{i \neq j < N} \text{Fix}_j, & \text{if } \iota = c_i^{++}, \\ \text{Dec}_i \wedge \bigwedge_{i \neq j < N} \text{Fix}_j, & \text{if } \iota = c_i^{--}, \\ \square_1^+ (C_i^+ \rightarrow C_i^-) \wedge \bigwedge_{j < N} \text{Fix}_j, & \text{if } \iota = c_i^{??}. \end{cases}$$

Now we can encode runs that start with all-0 counters. For each $q \in Q$, we take a fresh predicate symbol S_q , and define φ_M to be the conjunction of counter and the following formulas:

$$\bigwedge_{i < N} \square_1^+ (\neg C_i^+ \wedge \neg C_i^-), \tag{6}$$

$$\square_1^+ \square_0^+ (S \leftrightarrow \bigvee_{q \in Q-H} (S_q \wedge \bigwedge_{q \neq q' \in Q} \neg S_{q'})), \tag{7}$$

$$\square_1^+ \square_0^+ \bigwedge_{q \in Q-H} \left[S_q \rightarrow \bigvee_{\langle \iota, q' \rangle \in I_q} (\text{Do}_\iota \wedge \square_1 (N \rightarrow \square_0 (S \rightarrow S_{q'}))) \right]. \tag{8}$$

The following lemma says that going forward along the diagonal staircase generated in Claim 3.5, we can force infinite recurrent runs of M :

Lemma 3.7. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid} \wedge \varphi_M \wedge \Box_0 \Diamond_0 \Box_1 (S \rightarrow S_{q_r})$. For all $m < \omega$ and $i < N$, let*

$$q_m := q, \quad \text{if } \mathfrak{M}, \langle m, y_m \rangle \models S_q, \quad c_i(m) := |\{w \in W : \mathfrak{M}, \langle m, w \rangle \models C_i^+ \wedge \neg C_i^-\}|.$$

Then $\langle \langle q_m, \mathbf{c}(m) \rangle : m < \omega \rangle$ is a well-defined infinite run of M starting with all-0 counters and visiting q_r infinitely often.

Proof. The sequence $\langle q_m : m < \omega \rangle$ is well-defined and contains q_r infinitely often by Claim 3.5(ii), (7) and $\Box_0 \Diamond_0 \Box_1 (S \rightarrow S_{q_r})$. We show by induction on m that for all $m < \omega$,

$$\langle \langle q_0, \mathbf{c}(0) \rangle, \dots, \langle q_m, \mathbf{c}(m) \rangle \rangle$$

is a run of M starting with all-0 counters. Indeed, $c_i(0) = 0$ for $i < N$ by (6). Now suppose the statement holds for some $m < \omega$. By the IH, $\mathfrak{M}, \langle m, y_m \rangle \models S_{q_m}$. We have $q_m \in Q - H$ by Claim 3.5(ii) and (7), and so by (8) there is $\langle \iota, q' \rangle \in I_{q_m}$ such that $\mathfrak{M}, \langle m, y_m \rangle \models \text{Do}_\iota \wedge \Box_1 (N \rightarrow \Box_0 (S \rightarrow S_{q'}))$. Then $\mathfrak{M}, \langle m+1, y_{m+1} \rangle \models S_{q'}$ by Claim 3.5. Now there are three cases, depending on the form of ι . If $\iota = c_i^{++}$ for some $i < N$, then $c_i(m+1) = c_i(m) + 1$ and $c_j(m+1) = c_j(m)$, for $j \neq i$, $j < N$, by Claim 3.6. The case of $\iota = c_i^{--}$ is similar. If $\iota = c_i^{??}$ for some $i < N$, then $\mathfrak{M}, \langle m, y_m \rangle \models \Box_1^+(C_i^+ \rightarrow C_i^-)$, and so $c_i(m) = 0$. Also, $c_j(m+1) = c_j(m)$ for all $j < N$ by Claim 3.6. Therefore, in all cases we have $\langle q_m, \mathbf{c}(m) \rangle \rightarrow^t \langle q', \mathbf{c}(m+1) \rangle$, as required. \square

On the other hand, suppose M has an infinite run $\langle \langle q_m, \mathbf{c}(m) \rangle : m < \omega \rangle$ starting with all-0 counters and visiting q_r infinitely often. We define a model $\mathfrak{M}^{rec} = \langle \langle \omega, < \rangle \times \langle \omega, \neq \rangle, \rho \rangle$ as follows. For all $q \in Q$, we let

$$\begin{aligned} \rho(S) &:= \{ \langle n, n \rangle : n < \omega \}, \\ \rho(S_q) &:= \{ \langle n, n \rangle : n < \omega, q_n = q \}, \\ \rho(N) &:= \{ \langle n, n+1 \rangle : n < \omega \}. \end{aligned}$$

Further, for all $i < N$, $n < \omega$, we define inductively the sets $\rho_n(C_i^+)$ and $\rho_n(C_i^-)$. We let $\rho_0(C_i^+) = \rho_0(C_i^-) := \emptyset$, and

$$\begin{aligned} \rho_{n+1}(C_i^+) &:= \begin{cases} \rho_n(C_i^+) \cup \{n\}, & \text{if } \iota_n = c_i^{++}, \\ \rho_n(C_i^+), & \text{otherwise.} \end{cases} \\ \rho_{n+1}(C_i^-) &:= \begin{cases} \rho_n(C_i^-) \cup \{ \min(\rho_n(C_i^+) - \rho_n(C_i^-)) \}, & \text{if } \iota_n = c_i^{--}, \\ \rho_n(C_i^-), & \text{otherwise.} \end{cases} \end{aligned}$$

Finally, for each $i < N$, we let

$$\rho(C_i^+) := \{ \langle m, n \rangle : n \in \rho_m(C_i^+) \}, \quad \rho(C_i^-) := \{ \langle m, n \rangle : n \in \rho_m(C_i^-) \}.$$

It is straightforward to check that $\mathfrak{M}^{rec}, \langle 0, 0 \rangle \models \text{grid} \wedge \varphi_M \wedge \Box_0 \Diamond_0 \Box_1 (S \rightarrow S_{q_r})$, showing that CM recurrence is reducible to $\langle \omega, < \rangle \times \mathcal{C}_{diff}$ -satisfiability.

As concerns the Σ_1^1 upper bound, it is not hard to see that $\langle \omega, < \rangle \times \mathcal{C}_{diff}$ -satisfiability of a bimodal formula ϕ is expressible by a Σ_1^1 -formula over ω in the first-order language having binary predicate symbols $<$ and P^+ , for each propositional variable P in ϕ . This completes the proof of Theorem 3.1.

Next, we prove the lower bound of Theorem 3.2 by reducing the ‘CM reachability’ problem to $\mathcal{C}_{lin}^{fin} \times \mathcal{C}_{diff}$ -satisfiability. Let \mathfrak{M} be a model based on the product of some finite linear order $\langle T, < \rangle$ and some difference frame $\langle W, \neq \rangle$. We may assume that $T = |T| < \omega$. We encode counter machine runs in \mathfrak{M} like we did in the proof of Theorem 3.1, but of course this time only finite runs

are possible. We introduce a fresh propositional variable end , and let grid_{fin} be the conjunction of (3), (4) and the following version of (5):

$$\Box_0^+ \Box_1 (\mathbf{N} \wedge \neg \text{end} \rightarrow (\Diamond_0 S \wedge \Box_0 \Box_0 \neg S)). \quad (9)$$

The following finitary version of Claim 3.5 can be proved by a straightforward induction on m :

CLAIM 3.8. *Suppose $\mathfrak{M}, \langle 0, r \rangle \models \text{grid}_{fin}$. Then there exist some $0 < E \leq T$ and a sequence $\langle y_m \in W : m \leq E \rangle$ of points such that for all $m \leq E$,*

- (i) $y_0 = r$ and for all $n < m$, $y_m \neq y_n$,
- (ii) if $m < E$ then $\mathfrak{M}, \langle m, y_m \rangle \models S$,
- (iii) if $0 < m < E$ then $\mathfrak{M}, \langle m-1, y_m \rangle \models \mathbf{N}$,
- (iv) $\mathfrak{M}, \langle E-1, y_E \rangle \models \text{end}$, and if $0 < m < E-1$ then $\mathfrak{M}, \langle m-1, y_m \rangle \models \neg \text{end}$.

The proof of the following lemma is similar to that of Lemma 3.7:

Lemma 3.9. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid}_{fin} \wedge \varphi_M \wedge \Box_0^+ \Box_1 (\mathbf{N} \wedge \text{end} \rightarrow \Box_1 (S \rightarrow S_{q_r}))$. For all $m < E$ and $i < N$, let*

$$q_m := q, \quad \text{if } \mathfrak{M}, \langle m, y_m \rangle \models S_q, \quad c_i(m) := |\{w \in W : \mathfrak{M}, \langle m, w \rangle \models C_i^+ \wedge \neg C_i^-\}|.$$

Then $\langle \langle q_m, \mathbf{c}(m) \rangle : m < E \rangle$ is a well-defined run of M starting with all-0 counters and reaching q_r .

On the other hand, suppose M has a run $\langle \langle q_n, \mathbf{c}(n) \rangle : n < T \rangle$ for some $T < \omega$ such that it starts with all-0 counters and $q_{T-1} = q_r$. Take the model \mathfrak{M}^{rec} defined in the proof of Theorem 3.1 above. Let \mathfrak{M}^{fin} be its restriction to $\langle T, < \rangle \times \langle T+1, \neq \rangle$, and let

$$\rho(\text{end}) = \{\langle T-1, T \rangle\}.$$

Then it is straightforward to check that

$$\mathfrak{M}^{fin}, \langle 0, 0 \rangle \models \text{grid}_{fin} \wedge \varphi_M \wedge \Box_0^+ \Box_1 (\mathbf{N} \wedge \text{end} \rightarrow \Box_1 (S \rightarrow S_{q_r})),$$

completing the proof of the lower bound in Theorem 3.2.

As concerns the upper bound, recursively enumerability follows from the fact that $\mathcal{C}_{lin}^{fin} \times \mathcal{C}_{diff}$ -satisfiability has the ‘finite product model’ property:

CLAIM 3.10. *For any formula ϕ , if ϕ is $\mathcal{C}_{lin}^{fin} \times \mathcal{C}_{diff}$ -satisfiable, then ϕ is $\mathcal{C}_{lin}^{fin} \times \mathcal{C}_{diff}^{fin}$ -satisfiable.*

Proof. Suppose $\mathfrak{M}, \langle r_0, r_1 \rangle \models \phi$ for some model \mathfrak{M} based on the product of a finite linear order $\langle T, < \rangle$ and a (possibly infinite) difference frame $\langle W, \neq \rangle$. We may assume that $T = |T| < \omega$ and $r_0 = 0$. For all $n < T$, $X \subseteq W$, we define $cl_n(X)$ as the smallest set Y such that $X \subseteq Y \subseteq W$ and having the following property: If $x \in Y$ and $\mathfrak{M}, \langle n, x \rangle \models \Diamond_1 \psi$ for some $\psi \in \text{sub } \phi$, then there is $y \in Y$ such that $y \neq x$ and $\mathfrak{M}, \langle n, y \rangle \models \psi$. It is not hard to see that if X is finite then $cl_n(X)$ is finite as well. In fact, $|cl_n(X)| \leq |X| + 2|\text{sub } \phi|$. Now let $W_0 := cl_0(\{0\})$ and for $0 < n < T$ let $W_n := cl_n(W_{n-1})$. Let \mathfrak{M}' be the restriction of \mathfrak{M} to the product frame $\langle T, < \rangle \times \langle W_{T-1}, \neq \rangle$. An easy induction shows that for all $\psi \in \text{sub } \phi$, $n < T$, $w \in W_{T-1}$, we have $\mathfrak{M}, \langle n, w \rangle \models \psi$ iff $\mathfrak{M}', \langle n, w \rangle \models \psi$. \square

4 Undecidable bimodal logics with a ‘linear’ component

In this section we show that further combinations of weak order and pseudo-equivalence relations are undecidable. First, in Subsections 4.1 and 4.2 we show how to represent counter machine runs in ‘abstract’, not necessarily product frames for commutators. Then in Subsection 4.3 we extend our techniques to cover dense linear timelines. In order to obtain tighter control over the grid-structure, in all these cases we generate both the diagonal staircase and counter machine runs going *backward*, so the used formulas force infinite *rooted descending* chains in linear orders.

It is not clear, however, whether this change is always necessary, in other words, where exactly the limits of the ‘forward going’ technique are. In particular, it would be interesting to know whether the ‘infinite ascending chain’ analogues of the general Theorems 4.1 and 4.16 below hold.

4.1 Between commutators and products

In the following theorem we do not require the bimodal logic L to be Kripke complete:

Theorem 4.1. *Let L be any bimodal logic such that*

- L contains **[K4.3, Diff]**, and
- $\langle \omega + 1, > \rangle \times \langle \omega, \neq \rangle$ is a frame for L .

Then L -satisfiability is undecidable.

Corollary 4.2. *Both **[K4.3, Diff]** and **K4.3** \times **Diff** are undecidable.*

Note that Theorem 4.1 is much more general than Corollary 4.2, as not only **[K4.3, Diff]** \subsetneq **K4.3** \times **Diff**, but there are infinitely many different logics between them [20].

As a consequence of Theorems 2.1, 4.1 and Prop. 2.3 we also obtain:

Corollary 4.3. *FOLTL $^\neq$ -satisfiability is undecidable but co-r.e. in constant domain models over the class of all linear orders.*

We prove Theorem 4.1 by reducing ‘CM non-termination’ to L -satisfiability. To this end, fix some model \mathfrak{M} such that $\mathfrak{M} \models L$ and \mathfrak{M} is based on some 2-frame $\mathfrak{F} = \langle W, R_0, R_1 \rangle$. As by our assumption L -satisfiability of a formula implies its **[K4.3, Diff]**-satisfiability, by (2) we may assume that R_0 is transitive and weakly connected, R_1 is symmetric and pseudo-transitive, and R_0, R_1 commute. We begin with forcing a *unique* infinite diagonal staircase *backward*. Let grid^{bw} be the conjunction of the following formulas:

$$\Diamond_0(S \wedge \Box_0 \perp), \quad (10)$$

$$\Box_1^+ \Diamond_0 N, \quad (11)$$

$$\Box_1^+ \Box_0(N \rightarrow (\Box_1 \neg N \wedge \Diamond_1 S)), \quad (12)$$

$$\Box_1^+ \Box_0(N \rightarrow (\Diamond_0 S \wedge \Box_0 \Box_0 \neg S)), \quad (13)$$

$$\Box_1^+ \Box_0(S \rightarrow (\Box_0 \neg S \wedge \Box_1 \neg S)). \quad (14)$$

We will show, via a series of claims, that grid^{bw} forces not only a unique diagonal staircase, but also a unique ‘half-grid’ in \mathfrak{M} . To this end, for all $x \in W$, we define the *horizontal rank* of x by taking

$$hr(x) := \begin{cases} m, & \text{if the length of the longest } R_0\text{-path starting at } x \text{ is } m < \omega, \\ \omega, & \text{otherwise.} \end{cases}$$

CLAIM 4.4. *Suppose $\mathfrak{M}, r \models \text{grid}^{bw}$. Then there exist infinite sequences $\langle y_m : m < \omega \rangle$, $\langle u_m : m < \omega \rangle$, and $\langle v_m : m < \omega \rangle$ of points in W such that, for every $m < \omega$,*

- (i) $y_m = r$ or rR_1y_m , and $y_mR_0v_mR_0u_m$,

- (ii) if $m > 0$ then $v_{m-1}R_1u_m$,
- (iii) $\mathfrak{M}, u_m \models S$ and $hr(u_m) = m$,
- (iv) $\mathfrak{M}, v_m \models N$ and $hr(v_m) = m + 1$.

Proof. By induction on m . To begin with, let $y_0 = r$. By (10), there is u_0 such that $y_0R_0u_0$, $\mathfrak{M}, u_0 \models S$ and $hr(u_0) = 0$. By (11), there is v_0 such that $y_0R_0v_0$ and $\mathfrak{M}, v_0 \models N$. By (13), (14) and the weak connectedness of R_0 , we have that $v_0R_0u_0$, there is no x with $v_0R_0xR_0u_0$, and $hr(v_0) = 1$.

Now suppose inductively that for some $m < \omega$ we have y_k, u_k, v_k , for all $k \leq m$ as required. By the IH and (12), there is u_{m+1} such that $v_mR_1u_{m+1}$ and $\mathfrak{M}, u_{m+1} \models S$. As $hr(v_m) = m + 1$ by the IH, we have $hr(u_{m+1}) = m + 1$ by the commutativity of R_0 and R_1 . As $y_mR_0v_m$ by the IH, again by commutativity there is y_{m+1} such that $y_mR_1y_{m+1}R_0u_{m+1}$. As either $r = y_m$ or rR_1y_m by the IH and R_1 is pseudo-transitive, we have that either $r = y_{m+1}$ or rR_1y_{m+1} . So by (11), there is v_{m+1} such that $y_{m+1}R_0v_{m+1}$ and $\mathfrak{M}, v_{m+1} \models N$. By (13), (14) and the weak connectedness of R_0 , we have that $v_{m+1}R_0u_{m+1}$, there is no x with $v_{m+1}R_0xR_0u_{m+1}$, and $hr(v_{m+1}) = hr(u_{m+1}) + 1 = m + 2$ as required. \square

For each $m < \omega$, let $Column_m := \{u_m\} \cup \{x \in W : xR_1u_m\}$. The following claim is a straightforward consequence of Claim 4.4(iii), and the commutativity of R_0 and R_1 :

CLAIM 4.5. *For all $m < \omega$ and all $x \in Column_m$, $hr(x) = m$.*

Next, we define the half-grid points and prove some of their properties:

CLAIM 4.6. *Suppose that $\mathfrak{M}, r \models \text{grid}^{bw}$. Then for every pair $\langle m, n \rangle$ with $n < m < \omega$, there exists $x_{m,n} \in Column_m$ such that*

- (i) $x_{m,m-1} = v_{m-1}$, and if $n < m - 1$ then $x_{m,n}R_0x_{m-1,n}$,
- (ii) if $n < m - 1$ then there is no x with $x_{m,n}R_0xR_0x_{m-1,n}$.

Moreover, the $x_{m,n}$ are such that

- (iii) for all $x \in Column_m$, xR_0u_n iff $x = x_{m,n}$,
- (iv) $x_{m,n} \neq x_{m,n'}$ whenever $n \neq n'$.

Proof. First, by using Claim 4.4 throughout, we define some $x_{m,n} \in Column_m$ by induction on m satisfying (i) and (ii). To begin with, let $x_{1,0} = v_1$. Now suppose that $x_{m,n}$ satisfying (i) and (ii) have been defined for all $n < m$ for some $0 < m < \omega$. Take any $n < m + 1$. If $n = m$, then let $x_{m+1,m} = v_m$. If $n < m$ then $v_mR_0u_mR_1x_{m,n}$ by the IH. So by commutativity, there is $x_{m+1,n}$ such that $v_mR_1x_{m+1,n}R_0x_{m,n}$. As $u_{m+1}R_1v_m$, we have $x_{m+1,n} \in Column_{m+1}$ by the pseudo-transitivity of R_1 . Further, it follows from Claim 4.5 that there is no x with $x_{m+1,n}R_0xR_0x_{m,n}$.

Next, we show that the $x_{m,n}$ defined above satisfy (iii) and (iv). As $v_nR_0u_n$ by Claim 4.4(i), and $x_{m,n}R_0v_n$ by (i), we have $x_{m,n}R_0u_n$ by the transitivity of R_0 . For (iii): Let $x \in Column_m$ be such that xR_0u_n , and suppose that $x \neq x_{m,n}$. Then $xR_1x_{m,n}$, and so by commutativity, there is z with $x_{m,n}R_0zR_1u_n$. As R_0 is weakly connected and $hr(u_n) = hr(z)$ by Claim 4.5, we have $u_n = z$, and so $u_nR_1u_n$ follows. As $\mathfrak{M}, u_n \models S$ by Claim 4.4(iii), this contradicts (14), proving $x = x_{m,n}$. For (iv): Suppose, for contradiction, that $x_{m,n} = x_{m,n'}$ for some $n \neq n'$. By Claim 4.4(iii), $hr(u_n) = n \neq n' = hr(u_{n'})$, and so $u_n \neq u_{n'}$. As $x_{m,n}R_0u_n$ and $x_{m,n}R_0u_{n'}$, by the weak connectedness of R_0 , either $u_nR_0u_{n'}$ or $u_{n'}R_0u_n$. As $\mathfrak{M}, u_n \models S$ and $\mathfrak{M}, u_{n'} \models S$ by Claim 4.4(iii), this contradicts (14). \square

The following claim shows that we can in fact ‘single out’ the half-grid points in the columns by formulas:

CLAIM 4.7. *Suppose that $\mathfrak{M}, r \models \text{grid}^{bw}$. Then for all $m < \omega$ and all $x \in Column_m$,*

- (i) if $\mathfrak{M}, x \models \mathbf{N}$ then $m > 0$ and $x = v_{m-1} = x_{m,m-1}$,
- (ii) if $\mathfrak{M}, x \models \Diamond_0 \mathbf{N}$ then $m > 1$ and $x = x_{m,n}$ for some $0 < n < m - 1$.

Proof. Item (i) follows from Claim 4.4(iv) and (12). For (ii): Suppose that $\mathfrak{M}, x \models \Diamond_0 \mathbf{N}$ for some $x \in \text{Column}_m$. Then there is y such that xR_0y and $\mathfrak{M}, y \models \mathbf{N}$. By Claim 4.5, $hr(x) = m$, and so $hr(y) = n$ for some $n < m$. First, we claim that $x \neq x_{m,n}$. Indeed, suppose that $x = x_{m,n}$. Then by Claim 4.6, either $x = v_n$ or xR_0v_n . If $x = v_n$ then $\mathfrak{M}, x \models \mathbf{N}$ by Claim 4.4, contradicting (13). As $hr(v_n) = n + 1 > n = hr(y)$, $v_n \neq y$, the weak connectedness of R_0 and xR_0v_n imply that v_nR_0y , contradicting (13) again, and proving that $x \neq x_{m,n}$.

So we have $xR_1x_{m,n}$. By Claim 4.6, $x_{m,n}R_0u_n$. So by commutativity there is z such that $xR_0zR_1u_n$. Thus, $z \in \text{Column}_n$ and so $hr(z) = n$ by Claim 4.5. Then $y = z$ follows by the weak connectedness of R_0 , and so $y \in \text{Column}_n$. Thus, we have $n > 0$ and $y = v_{n-1}$ by (i). Therefore, $m > 1$, and $xR_0v_{n-1}R_0u_{n-1}$ by Claim 4.4. So $x = x_{m,n-1}$ follows by Claim 4.6(iii). \square

Given a counter machine M , we now encode runs that start with all-0 counters by going backward along the created diagonal staircase. For each counter $i < N$, we take a fresh propositional variable C_i . At each moment n of time, the content of counter c_i at step n of a run is represented by those points in Column_n where C_i holds. We also force these points only to be among the half-grid points $x_{m,n}$. We can achieve these by the following formula:

$$\begin{aligned} \text{counter}^{bw} &:= \Box_1^+ \Box_0 \bigwedge_{i < N} (C_i \rightarrow (\mathbf{N} \vee \text{All}C_i)), \quad \text{where} \\ \text{All}C_i &:= \Diamond_0 \mathbf{N} \wedge \Box_0 (\mathbf{N} \vee \Diamond_0 \mathbf{N} \rightarrow C_i). \end{aligned} \tag{15}$$

CLAIM 4.8. *Suppose that $\mathfrak{M}, r \models \text{grid}^{bw} \wedge \text{counter}^{bw}$. Then for all $m < \omega$, $i < N$,*

$$|\{x \in \text{Column}_{m+1} : \mathfrak{M}, x \models \text{All}C_i\}| = |\{x \in \text{Column}_m : \mathfrak{M}, x \models C_i\}|.$$

Proof. As $\Diamond_0 \mathbf{N}$ is a conjunct of $\text{All}C_i$, by Claims 4.6(iv), 4.7 and counter^{bw} , we have

$$\begin{aligned} |\{x \in \text{Column}_{m+1} : \mathfrak{M}, x \models \text{All}C_i\}| &= |\{n : n < m \text{ and } \mathfrak{M}, x_{m+1,n} \models \text{All}C_i\}|, \text{ and} \\ |\{x \in \text{Column}_m : \mathfrak{M}, x \models C_i\}| &= |\{n : n < m \text{ and } \mathfrak{M}, x_{m,n} \models C_i\}|. \end{aligned}$$

So it is enough to show that the two sets on the right hand sides are equal. To this end, suppose first that $n < m$ is such that $\mathfrak{M}, x_{m+1,n} \models \text{All}C_i$. As $x_{m+1,n}R_0x_{m,n}$ by Claim 4.6(i), and $\mathfrak{M}, x_{m,n} \models \mathbf{N} \vee \Diamond_0 \mathbf{N}$ by Claims 4.4(iv) and 4.6(i), we obtain that $\mathfrak{M}, x_{m,n} \models C_i$.

Conversely, suppose that $\mathfrak{M}, x_{m,n} \models C_i$ for some $n < m$. As $n < m$, by Claims 4.4(iv) and 4.6(i), we have $\mathfrak{M}, x_{m+1,n} \models \Diamond_0 \mathbf{N}$. Now let x be such that $x_{m+1,n}R_0x$ and $\mathfrak{M}, x \models \mathbf{N} \vee \Diamond_0 \mathbf{N}$. By Claim 4.7 and the weak connectedness of R_0 , either $x = x_{m,n}$ or $x_{m,n}R_0x$. In the former case, $\mathfrak{M}, x \models C_i$ by assumption. If $x_{m,n}R_0x$ then $\mathfrak{M}, x_{m,n} \models \neg \mathbf{N}$ by (13). Therefore, $\mathfrak{M}, x_{m,n} \models \text{All}C_i$ by (15), and so $\mathfrak{M}, x_{m,n} \models \Box_0 (\mathbf{N} \vee \Diamond_0 \mathbf{N} \rightarrow C_i)$. Thus, we have $\mathfrak{M}, x \models C_i$ in this case as well, and so $\mathfrak{M}, x_{m+1,n} \models \Box_0 (\mathbf{N} \vee \Diamond_0 \mathbf{N} \rightarrow C_i)$ as required. \square

Now, for each $i < N$, the following formulas simulate the possible changes that may happen in the counters when stepping backward, and also ensure that each ‘vertical coordinate’ is used only once in the counting:

$$\text{Fix}_i^{bw} := \Box_1^+ (C_i \leftrightarrow \text{All}C_i), \tag{16}$$

$$\text{Inc}_i^{bw} := \Box_1^+ (C_i \leftrightarrow (\mathbf{N} \vee \text{All}C_i)), \tag{17}$$

$$\text{Dec}_i^{bw} := \Box_1^+ (C_i \rightarrow \text{All}C_i) \wedge \Diamond_1^1 (\neg C_i \wedge \text{All}C_i). \tag{18}$$

The following analogue of Claim 3.6 is a straightforward consequence of Claim 4.8:

CLAIM 4.9. Suppose that $\mathfrak{M}, r \models \text{grid}^{bw} \wedge \text{counter}^{bw}$ and let, for all $m < \omega$, $i < N$, $c_i(m) := |\{x \in \text{Column}_m : \mathfrak{M}, x \models C_i\}|$. Then

$$c_i(m+1) = \begin{cases} c_i(m), & \text{if } \mathfrak{M}, u_{m+1} \models \text{Fix}_i^{bw}, \\ c_i(m) + 1, & \text{if } \mathfrak{M}, u_{m+1} \models \text{Inc}_i^{bw}, \\ c_i(m) - 1, & \text{if } \mathfrak{M}, u_{m+1} \models \text{Dec}_i^{bw}. \end{cases}$$

Next, we encode the various counter machine instructions, acting backward. For each $\iota \in \text{Op}_C$, we define the formula Do_ι^{bw} by taking

$$\text{Do}_\iota^{bw} := \begin{cases} \text{Inc}_i^{bw} \wedge \bigwedge_{i \neq j < N} \text{Fix}_j^{bw}, & \text{if } \iota = c_i^{++}, \\ \text{Dec}_i^{bw} \wedge \bigwedge_{i \neq j < N} \text{Fix}_j^{bw}, & \text{if } \iota = c_i^{--}, \\ \Box_1^+ \neg C_i \wedge \bigwedge_{j < N} \text{Fix}_j^{bw}, & \text{if } \iota = c_i^{??}. \end{cases}$$

Finally, we encode runs that start with all-0 counters. For each $\iota \in \text{Op}_C$, we introduce a propositional variable \mathbf{l}_ι , and define φ_M^{bw} to be the conjunction of counter^{bw} and the following formulas:

$$\Box_1^+ \Box_0 (S \leftrightarrow \bigvee_{q \in Q-H} (S_q \wedge \bigwedge_{q \neq q' \in Q} \neg S_{q'})), \quad (19)$$

$$\Box_1 \Box_0 \bigwedge_{q \in Q-H} [(S \wedge \Diamond_1 (N \wedge \Diamond_0 S_q)) \rightarrow \bigvee_{\langle \iota, q' \rangle \in I_q} (\mathbf{l}_\iota \wedge S_{q'})], \quad (20)$$

$$\Box_1 \Box_0 \bigwedge_{\iota \in \text{Op}_C} (\mathbf{l}_\iota \rightarrow \text{Do}_\iota^{bw}). \quad (21)$$

The following analogue of Lemma 3.7 says that going backward along the diagonal staircase generated in Claim 4.4, we can force infinite runs of M :

Lemma 4.10. Suppose that $\mathfrak{M}, r \models \text{grid}^{bw} \wedge \varphi_M^{bw}$, and for all $m < \omega$ and $i < N$, let

$$q_m := q, \text{ if } \mathfrak{M}, u_m \models S_q, \quad c_i(m) := |\{x \in \text{Column}_m : \mathfrak{M}, x \models C_i\}|, \quad \sigma_m := \langle q_m, \mathbf{c}(m) \rangle.$$

Then $\langle \sigma_m : m < \omega \rangle$ is a well-defined infinite run of M starting with all-0 counters.

Proof. The sequence $\langle q_m : m < \omega \rangle$ is well-defined by Claim 4.4(iii) and (19). We show by induction on m that for all $m < \omega$, $\langle \sigma_0, \dots, \sigma_m \rangle$ is a run of M starting with all-0 counters. Indeed, $c_i(0) = 0$ for $i < N$ by (15) and Claim 4.7. Now suppose the statement holds for some $m < \omega$. By Claim 4.4, $\mathfrak{M}, u_{m+1} \models S \wedge \Diamond_1 (N \wedge \Diamond_0 S_{q_m})$. By (19) we have $q_m \in Q - H$, and so by (20) there is $\langle \iota, q_{m+1} \rangle \in I_{q_m}$ such that $\mathfrak{M}, u_{m+1} \models \mathbf{l}_\iota \wedge S_{q_{m+1}}$. Therefore, so $\mathfrak{M}, u_{m+1} \models \text{Do}_\iota^{bw}$ by (21). It follows from Claim 4.9 that $\sigma_m \rightarrow^\iota \sigma_{m+1}$ as required. \square

On the other hand, suppose that M has an infinite run $\langle \sigma_n : n < \omega \rangle$ starting with all-0 counters such that $\sigma_n = \langle q_n, \mathbf{c}_n \rangle$ and $\sigma_n \rightarrow^{\iota_n} \sigma_{n+1}$, for $n < \omega$. We define a model

$$\mathfrak{M}^\infty = \langle \langle \omega + 1, > \rangle \times \langle \omega, \neq \rangle, \mu \rangle$$

as follows. For all $q \in Q$ and $\iota \in \text{Op}_C$, we let

$$\mu(S) := \{ \langle n, n \rangle : n < \omega \}, \quad (22)$$

$$\mu(S_q) := \{ \langle n, n \rangle : n < \omega, q_n = q \}, \quad (23)$$

$$\mu(N) := \{ \langle n + 1, n \rangle : n < \omega \}, \quad (24)$$

$$\mu(\mathbf{l}_\iota) := \{ \langle n, n \rangle : n < \omega, \iota = \iota_n \}. \quad (25)$$

Further, for all $i < N$, $n < \omega$, we define inductively the sets $\mu_n(\mathbf{C}_i)$. We let $\mu_0(\mathbf{C}_i) := \emptyset$, and

$$\mu_{n+1}(\mathbf{C}_i) := \begin{cases} \mu_n(\mathbf{C}_i) \cup \{n\}, & \text{if } \iota_n = c_i^{++}, \\ \mu_n(\mathbf{C}_i) - \{\min(\mu_n(\mathbf{C}_i))\}, & \text{if } \iota_n = c_i^{--}, \\ \mu_n(\mathbf{C}_i), & \text{otherwise.} \end{cases} \quad (26)$$

Finally, for each $i < N$, we let

$$\mu(\mathbf{C}_i) := \{ \langle m, n \rangle : m < \omega, n \in \mu_m(\mathbf{C}_i) \}. \quad (27)$$

It is straightforward to check that $\mathfrak{M}^\infty, \langle \omega, 0 \rangle \models \text{grid}^{bw} \wedge \varphi_M^{bw}$, showing that CM non-termination can be reduced to L -satisfiability. This completes the proof of Theorem 4.1.

4.2 Modally discrete weak orders with infinite descending chains

In some cases, we can have stronger lower bounds than in Theorem 4.1. We call a frame $\langle W, R \rangle$ *modally discrete* if it satisfies the following aspect of discreteness: there are no points $x_0, x_1, \dots, x_n, \dots, x_\infty$ in W such that $x_0 R x_1 R x_2 R \dots R x_n R \dots$, $x_i \neq x_{i+1}$, $x_i R x_\infty$ and $x_\infty \neg R x_i$, for all $i < \omega$. We denote by **DisK4.3** the logic of all modally discrete weak orders. Several well-known ‘linear’ modal logics are extensions of **DisK4.3**, for example, **Log** $\langle \omega, < \rangle$ and **GL.3** (the logic of all Noetherian⁶ irreflexive linear orders). Unlike ‘real’ discreteness, modal discreteness can be captured by modal formulas, and each of these logics is finitely axiomatisable [42, 11]. Also, note that for $L \in \{\mathbf{DisK4.3}, \mathbf{Log}\langle \omega, < \rangle, \mathbf{GL.3}\}$, either $\langle \omega + 1, > \rangle$ or $\langle \{\infty\} \cup \mathbb{Z}, > \rangle$ is a frame for L (here \mathbb{Z} denotes the set of all integers).

Theorem 4.11. *Let \mathcal{C} be any class of frames for $[\mathbf{DisK4.3}, \mathbf{Diff}]$ such that either $\langle \omega + 1, > \rangle \times \langle \omega, \neq \rangle$ or $\langle \{\infty\} \cup \mathbb{Z}, > \rangle \times \langle \omega, \neq \rangle$ belongs to \mathcal{C} . Then \mathcal{C} -satisfiability is Σ_1^1 -hard.*

Corollary 4.12. *Let L_1 be any logic from the list*

$$\mathbf{Log}\langle \omega, < \rangle, \mathbf{GL.3}, \mathbf{DisK4.3}.$$

Then, for any Kripke complete bimodal logic L in the interval

$$[L_1, \mathbf{Diff}] \subseteq L \subseteq L_1 \times \mathbf{Diff},$$

L -satisfiability is Σ_1^1 -hard.

We also obtain the following interesting corollary. As $[L_0, L_1]$ -satisfiability is clearly co-r.e whenever both L_0 and L_1 are finitely axiomatisable, Corollary 4.12 yields new examples of *Kripke incomplete* commutators of Kripke complete and finitely axiomatisable logics:

Corollary 4.13. *Let L_1 be like in Corollary 4.12. Then the commutator $[L_1, \mathbf{Diff}]$ is Kripke incomplete.*

Note that it is not known whether any of the commutators $[L_1, \mathbf{S5}]$ is decidable or Kripke complete, whenever L_1 is one of the logics in Corollary 4.12.

We prove Theorem 4.11 by reducing the ‘CM recurrence’ problem to \mathcal{C} -satisfiability. Let \mathfrak{M} be a model over some 2-frame $\mathfrak{F} = \langle W, R_0, R_1 \rangle$ in \mathcal{C} . As **DisK4.3** \supseteq **K4.3**, \mathfrak{F} is a frame for **[K4.3, Diff]**. So by (2) we may assume that R_0 is a modally discrete weak order, R_1 is symmetric and pseudo-transitive, and R_0, R_1 commute. We will encode counter machine runs in \mathfrak{M} ‘going backward’, like we did in the proof of Theorem 4.1, with the help of the formulas grid^{bw} and φ_M^{bw} . This time we use some additional machinery ensuring recurrence. To this end, we introduce

⁶ $\langle W, R \rangle$ is *Noetherian* if it contains no infinite ascending chains $x_0 R x_1 R x_2 R \dots$ where $x_i \neq x_{i+1}$.

two fresh propositional variables R and Q , and define the formula rec^{bw} as the conjunction of the following formulas:

$$\Box_1^+ \Box_0 (S \rightarrow \Diamond_1 R), \quad (28)$$

$$\Box_1^+ \Box_0 (R \rightarrow \Box_0 \neg S), \quad (29)$$

$$\Box_0 (\Diamond_1 S \rightarrow \Diamond_1 N), \quad (30)$$

$$\Box_1 \Box_0 \left[S \rightarrow \left(Q \leftrightarrow \Box_1 (N \rightarrow \Box_0 (S \rightarrow \neg Q)) \right) \right], \quad (31)$$

$$\Box_1^+ \Box_0 (S \wedge \Diamond_0 R \rightarrow S_{q_r}), \quad (32)$$

where q_r is the state of counter machine M we will force to recur. In the following claim and its proof we use the notation introduced in Claims 4.4–4.6:

CLAIM 4.14. *Suppose that $\mathfrak{M}, r \models \text{grid}^{bw} \wedge \text{rec}^{bw}$. Then there are infinitely many m such that $\mathfrak{M}, u_m \models S_{q_r}$.*

Proof. We show that for every $m < \omega$ there is $k_m > m$ with $\mathfrak{M}, u_{k_m} \models S_{q_r}$. Fix any $m < \omega$. By Claim 4.4(iii) and (28), there is w^* such that $u_m R_1 w^*$ and $\mathfrak{M}, w^* \models R$. We claim that

$$\text{there is } k < \omega \text{ such that } u_k R_0 w^*. \quad (33)$$

Indeed, suppose for contradiction that (33) does not hold. We define by induction a sequence $\langle x_n : n < \omega \rangle$ of points such that, for all $n < \omega$,

$$r R_0 x_n, \quad (34)$$

$$x_n \notin \text{Column}_k \text{ for any } k < \omega, \quad (35)$$

$$\mathfrak{M}, x_n \models \Diamond_1 S, \quad (36)$$

$$\text{if } n > 0 \text{ then } x_{n-1} R_0 x_n \text{ and } x_{n-1} \neq x_n. \quad (37)$$

To begin with, by commutativity of R_0 and R_1 , we have some y with $r R_1 y R_0 w^*$. So by (11), there is b_0 such that $y R_0 b_0$ and $\mathfrak{M}, b_0 \models N$. By (13), there is a_0 such that $b_0 R_0 a_0$, there is no b with $b_0 R_0 b R_0 a_0$ and $\mathfrak{M}, a_0 \models S$. By commutativity, there is x_0 such that $r R_0 x_0 R_1 a_0$, and so $\mathfrak{M}, x_0 \models \Diamond_1 S$. By Claim 4.4(iii), (14), (29) and the weak connectedness of R_0 , we have $a_0 R_0 w^*$. Therefore, $a_0 \neq u_k$ for any $k < \omega$ by our indirect assumption, and so $a_0 \notin \text{Column}_k$ for any $k < \omega$ by (14). As $x_0 R_1 a_0$, it follows that $x_0 \notin \text{Column}_k$ for any $k < \omega$.

Now suppose inductively that we have $\langle x_i : i \leq n \rangle$ satisfying (34)–(37) for some $n < \omega$. By (36) of the IH and (30), there is b_{n+1} such that $x_n R_1 b_{n+1}$ and $\mathfrak{M}, b_{n+1} \models N$. By (13), there is a_{n+1} such that $b_{n+1} R_0 a_{n+1}$, there is no b with $b_{n+1} R_0 b R_0 a_{n+1}$ and $\mathfrak{M}, a_{n+1} \models S$. By commutativity, there is x_{n+1} such that $x_n R_0 x_{n+1} R_1 a_{n+1}$, and so $r R_0 x_{n+1}$ and $\mathfrak{M}, x_{n+1} \models \Diamond_1 S$. We claim that

$$x_{n+1} \neq x_n. \quad (38)$$

Suppose for contradiction that $x_{n+1} = x_n$. Let a_n be such that $x_n R_1 a_n$ and $\mathfrak{M}, a_n \models S$. Then $a_n = a_{n+1}$ follows by (14). However, by (13), (14) and (31) we obtain that $a_n \neq a_{n+1}$. So we have a contradiction, proving (38). Finally, we claim that

$$x_{n+1} \notin \text{Column}_k \text{ for any } k < \omega. \quad (39)$$

Suppose not, that is, $x_{n+1} \in \text{Column}_k$ for some $k < \omega$. As $x_{n+1} R_1 a_{n+1}$, we also have that $a_{n+1} \in \text{Column}_k$. Then $hr(b_{n+1}) = k + 1$, by the weak connectedness of R_0 and Claim 4.5, and so $hr(x_n) = k + 1$ by $x_n R_1 b_{n+1}$ and commutativity. Take the grid-point $x_{k+1,0} \in \text{Column}_{k+1}$ defined in Claim 4.6. As $hr(x_{k+1,0}) = k + 1$ by Claim 4.5, we have $x_{k+1,0} = x_n$ by the weak connectedness of R_0 . But this contradicts (35) of the IH, proving (39).

So we have defined $\langle x_n : n < \omega \rangle$ satisfying (34)–(37). As $hr(u_0) = 0$ by Claim 4.4(iii), and $\mathfrak{M}, x_n \models \neg S$ by (14) and (36), by the weak connectedness of R_0 we obtain that $x_n R_0 u_0$ for every $n < \omega$. This contradicts the modal discreteness of R_0 , and so proves (33).

Now let k_m be such that $u_{k_m} R_0 w^*$. As $w^* \in \text{Column}_{k_m}$, $k_m > m$ follows from Claim 4.5. By Claim 4.4(iii) and (32), we have $\mathfrak{M}, u_{k_m} \models S_{q_r}$ as required. \square

Now the following lemma is a straightforward consequence of Lemma 4.10 and Claim 4.14:

Lemma 4.15. *Suppose that $\mathfrak{M}, r \models \text{grid}^{bw} \wedge \varphi_M^{bw} \wedge \text{rec}^{bw}$, and for all $m < \omega$, $i < N$, let*

$$q_m := q, \text{ if } \mathfrak{M}, u_m \models S_q, \quad c_i(m) := |\{x \in \text{Column}_m : \mathfrak{M}, x \models C_i\}|, \quad \sigma_m := \langle q_m, \mathbf{c}(m) \rangle.$$

Then $\langle \sigma_m : m < \omega \rangle$ is a well-defined run of M starting with all-0 counters and visiting q_r infinitely often.

On the other hand, suppose that M has run $\langle \langle q_n, \mathbf{c}(n) \rangle : n < \omega \rangle$ such that $\mathbf{c}(0) = 0$ and $q_{k_n} = q_r$ for an infinite sequence $\langle k_n : n < \omega \rangle$. Clearly, we may assume that $k_n > n$, for $n < \omega$. By assumption, $\mathfrak{F} \in \mathcal{C}$ for either $\mathfrak{F} = \langle \omega + 1, > \rangle \times \langle \omega, \neq \rangle$ or $\mathfrak{F} = \langle \{\infty\} \cup \mathbb{Z}, > \rangle \times \langle \omega, \neq \rangle$. Then the model \mathfrak{M}^∞ defined in (22)–(27) can be regarded as a model based on \mathfrak{F} , and we may add

$$\mu(Q) := \{ \langle n, n \rangle : n < \omega, n \text{ is odd} \}, \quad \mu(R) := \{ \langle n, k_n \rangle : n < \omega \}.$$

It is straightforward to check that $\mathfrak{M}^\infty, \langle \omega, 0 \rangle \models \text{grid}^{bw} \wedge \varphi_M^{bw} \wedge \text{rec}^{bw}$. So by Lemma 4.15, CM recurrence can be reduced to \mathcal{C} -satisfiability, proving Theorem 4.11.

4.3 Decreasing 2-frames based on dense weak orders

A weak order $\langle W, R \rangle$ is called *dense* if $\forall x, y (xRy \rightarrow \exists z xRzRy)$. Well-known examples of dense linear orders are $\langle \mathbb{Q}, < \rangle$ and $\langle \mathbb{R}, < \rangle$ of the *rationals* and the *reals*, respectively. Neither Theorem 3.1 nor Theorem 4.1 apply if the ‘horizontal component’ of a bimodal logic has only dense frames. In this section we cover some of these cases.

We say that a frame $\mathfrak{F} = \langle W, R \rangle$ *contains an $\langle \omega + 1, > \rangle$ -type chain*, if there are distinct points x_n , for $n \leq \omega$, in W such that $x_n R x_m$ iff $n > m$, for all $n, m \leq \omega$, $n \neq m$. Observe that this is less than saying that \mathfrak{F} has a subframe isomorphic to $\langle \omega + 1, > \rangle$, as for each n , $x_n R x_n$ might or might not hold. So \mathfrak{F} can be reflexive and/or dense, and still have this property. We have the following generalisation of Theorem 4.1 for classes of decreasing 2-frames:

Theorem 4.16. *Let \mathcal{C} be any class of weak orders such that $\mathfrak{F} \in \mathcal{C}$ for some \mathfrak{F} containing an $\langle \omega + 1, > \rangle$ -type chain. Then $\mathcal{C} \times {}^d\mathcal{C}_{\text{diff}}$ -satisfiability is undecidable.*

As a consequence of Theorem 4.16 and Props. 2.2, 2.3 we obtain:

Corollary 4.17. *FOLTL $^\neq$ -satisfiability is undecidable both in decreasing and in constant domain models over $\langle \mathbb{Q}, < \rangle$ and over $\langle \mathbb{R}, < \rangle$.*

Also, as a consequence of Theorems 2.1, 4.16 and Props. 2.2, 2.3 we have:

Corollary 4.18. *FOLTL $^\neq$ -satisfiability is undecidable but co-r.e. in decreasing domain models over the class of all linear orders.*

We prove Theorem 4.16 by reducing the ‘CM non-termination’ problem to $\mathcal{C} \times {}^d\mathcal{C}_{\text{diff}}$ -satisfiability. We intend to use something like the formula $\text{grid}^{bw} \wedge \varphi_M^{bw}$ defined in the proof of Theorem 4.1. The problem is that if $\langle W, R \rangle$ is reflexive and/or dense, then a formula of the form $\Diamond_0 S \wedge \Box_0 \Box_0 \neg S$ in conjunct (13) of grid^{bw} is clearly not satisfiable. In order to overcome this, we will apply a version of the well-known ‘tick trick’ (see e.g. [46, 39, 15]).

So let \mathfrak{M} be a model based on a decreasing 2-frame $\mathfrak{H}_{\mathfrak{F}, \mathfrak{G}}$ where $\mathfrak{F} = \langle W, R \rangle$ is a weak order, and for every $x \in W$, $\mathfrak{G}_x = \langle W_x, \neq \rangle$. We may assume that \mathfrak{F} is rooted with some r_0 as its root. We take a fresh propositional variable *Tick*, and define a new modal operator by setting, for every formula ψ ,

$$\begin{aligned} \blacklozenge_0 \psi &:= [\text{Tick} \wedge \Diamond_0 (\neg \text{Tick} \wedge (\psi \vee \Diamond_0 \psi))] \vee [\neg \text{Tick} \wedge \Diamond_0 (\text{Tick} \wedge (\psi \vee \Diamond_0 \psi))], \text{ and} \\ \blacksquare_0 \phi &:= \neg \blacklozenge_0 \neg \phi. \end{aligned}$$

Now suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models (40)$, where

$$\Box_1^+ \Box_0^+ (\text{Tick} \vee \Diamond_1 \text{Tick} \rightarrow (\text{Tick} \wedge \Box_1 \text{Tick})). \quad (40)$$

We define a new binary relation $R^{\mathfrak{M}}$ on W by taking, for all $x, y \in W$,

$$xR^{\mathfrak{M}}y \quad \text{iff} \quad \exists z \in W (xRz \text{ and } (z = y \text{ or } zRy) \text{ and } \forall u \in W_z (\mathfrak{M}, \langle x, u \rangle \models \text{Tick} \leftrightarrow \mathfrak{M}, \langle z, u \rangle \models \neg \text{Tick})).$$

Then it is not hard to check that $R^{\mathfrak{M}}$ is transitive, and \blacklozenge_0 behaves like a ‘horizontal’ modal diamond w.r.t. $R^{\mathfrak{M}}$ in \mathfrak{M} , that is, for all $x \in W$, $u \in W_x$,

$$\mathfrak{M}, \langle x, u \rangle \models \blacklozenge_0 \psi \quad \text{iff} \quad \exists y \in W (xR^{\mathfrak{M}}y, u \in W_y \text{ and } \mathfrak{M}, \langle y, u \rangle \models \psi).$$

However, $R^{\mathfrak{M}}$ is not necessarily weakly connected. We only have:

$$\forall x, y, z (xR^{\mathfrak{M}}y \wedge xR^{\mathfrak{M}}z \rightarrow (y \sim z \vee yR^{\mathfrak{M}}z \vee zR^{\mathfrak{M}}y)), \quad (41)$$

where

$$y \sim z \quad \text{iff} \quad \text{either } y = z \text{ or } (yRz \text{ and } y \neg R^{\mathfrak{M}}z) \text{ or } (zRy \text{ and } z \neg R^{\mathfrak{M}}y).$$

The relation \sim can be genuinely larger than equality. It is not hard to check (using that $\langle W, R \rangle$ is rooted) that \sim is an equivalence relation, and \sim -related points have the following properties:

$$\forall x, y, z (y \sim z \wedge xR^{\mathfrak{M}}y \rightarrow xR^{\mathfrak{M}}z), \quad (42)$$

$$\forall x, y, z (y \sim z \wedge yR^{\mathfrak{M}}x \rightarrow zR^{\mathfrak{M}}x). \quad (43)$$

We would like our propositional variables to behave ‘uniformly’ when interpreted at pairs with \sim -related first components (that is, along ‘horizontal intervals’). To achieve this, for a propositional variable P , let Interval_P denote conjunction of the following formulas:

$$\Box_1^+ \Box_0^+ (P \rightarrow \blacksquare_0 \neg P), \quad (44)$$

$$\Box_1^+ \Box_0^+ (\Diamond_0 P \wedge \blacksquare_0 \neg P \rightarrow P), \quad (45)$$

$$\Box_1^+ \Box_0^+ (P \wedge \neg \blacklozenge_0 \top \rightarrow \Box_0 P), \quad (46)$$

$$\Box_1^+ \Box_0^+ (P \wedge \blacklozenge_0 \top \rightarrow \blacklozenge_0 P'), \quad (47)$$

$$\Box_1^+ \Box_0^+ (P \rightarrow \Box_0 (\blacklozenge_0 P' \rightarrow P)), \quad (48)$$

where P' is a fresh propositional variable. We also introduce the following notation, for all $x \in W$, $y \in W_x$ and all formulas ϕ :

$$\mathfrak{M}, \langle I(x), y \rangle \models \phi \quad \text{iff} \quad \mathfrak{M}, \langle z, y \rangle \models \phi \text{ for all } z \text{ such that } z \sim x \text{ and } y \in W_z.$$

CLAIM 4.19. *Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models (40) \wedge \text{Interval}_P$. For all $x \in W$, $y \in W_x$, if $\mathfrak{M}, \langle x, y \rangle \models P$ then $\mathfrak{M}, \langle I(x), y \rangle \models P$.*

Proof. Suppose that $\mathfrak{M}, \langle x, y \rangle \models P$. Take some $z \sim x$ with $z \neq x$ and $y \in W_z$. Suppose first that zRx . As $\mathfrak{M}, \langle x, y \rangle \models \blacksquare_0 \neg P$ by (44), we have $\mathfrak{M}, \langle z, y \rangle \models \blacksquare_0 \neg P$ by (43). Therefore, $\mathfrak{M}, \langle z, y \rangle \models P$ by (45).

Now suppose that xRz . There are two cases: If $\mathfrak{M}, \langle x, y \rangle \models \neg \blacklozenge_0 \top$ then $\mathfrak{M}, \langle z, y \rangle \models P$ follows by (46). If $\mathfrak{M}, \langle x, y \rangle \models \blacklozenge_0 \top$ then $\mathfrak{M}, \langle x, y \rangle \models \blacklozenge_0 P'$ by (47). Thus, $\mathfrak{M}, \langle z, y \rangle \models \blacklozenge_0 P'$ by (43). So $\mathfrak{M}, \langle z, y \rangle \models P$ follows by (48). \square

Throughout, for any formula ϕ , we denote by ϕ^\bullet the formula obtained from ϕ by replacing each occurrence of \Diamond_0 with \blacklozenge_0 . Now all the necessary tools are ready for forcing a unique infinite diagonal staircase of intervals, going backward. In decreasing 2-frames this will also automatically give us an infinite half-grid. To this end, take the formula grid^{bw} defined in (10)–(14). We define a

new formula grid^* by modifying grid^{bw} as follows. First, replace the conjunct (10) by the slightly stronger

$$\Diamond_0(S \wedge \Box_1^+ \blacksquare_0 \perp), \quad (49)$$

then replace each remaining conjunct ϕ in grid^{bw} by ϕ^\bullet . Finally, add the conjuncts (40) and Interval_P , for $P \in \{N, S\}$. We then have the following analogue of Claims 4.4–4.6:

CLAIM 4.20. *Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models \text{grid}^*$. Then there exist infinite sequences $\langle x_m \in W : m < \omega \rangle$ and $\langle y_m \in W_{x_m} : m < \omega \rangle$ such that for all $m < \omega$,*

- (i) $y_m \neq y_n$, for all $n < m$,
- (ii) there is no x with $x_0 R^{\mathfrak{M}} x$, and if $m > 0$ then $x_m R^{\mathfrak{M}} x_{m-1}$, and there is no x such that $x_m R^{\mathfrak{M}} x R^{\mathfrak{M}} x_{m-1}$,
- (iii) $\mathfrak{M}, \langle I(x_m), y_m \rangle \models S$,
- (iv) if $m > 0$ then $\mathfrak{M}, \langle I(x_m), y_{m-1} \rangle \models N$.

Proof. By induction on m . To begin with, let $y_0 = r_1$. By (49), there is x_0 such that $r_0 R^{\mathfrak{M}} x_0$, $y_0 \in W_{x_0}$, $\mathfrak{M}, \langle x_0, y_0 \rangle \models S$ and

$$\mathfrak{M}, \langle x_0, y_0 \rangle \models \Box_1^+ \blacksquare_0 \perp. \quad (50)$$

By Intervals_S , we have $\mathfrak{M}, \langle I(x_0), y_0 \rangle \models S$.

Now suppose inductively that for some $m < \omega$ we have x_k, y_k , for all $k \leq m$ as required. By the IH, $y_m \in W_{x_m} \subseteq W_{r_0}$, so by (11) $^\bullet$, there is x_{m+1} such that $r_0 R^{\mathfrak{M}} x_{m+1}$, $y_m \in W_{x_{m+1}}$ and $\mathfrak{M}, \langle x_{m+1}, y_m \rangle \models N$. By (13) $^\bullet$, (14) $^\bullet$, (41) and (43), we have that $x_{m+1} R^{\mathfrak{M}} x_m$, and there is no x with $x_{m+1} R^{\mathfrak{M}} x R^{\mathfrak{M}} x_{m+1}$. By Interval_N , we have $\mathfrak{M}, \langle I(x_{m+1}), y_m \rangle \models N$. By (12) $^\bullet$, there is y_{m+1} such that $y_{m+1} \neq y_m$, $y_{m+1} \in W_{x_m}$ and $\mathfrak{M}, \langle x_{m+1}, y_{m+1} \rangle \models S$. By Intervals_S , we have $\mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models S$. Finally, we have $y_{m+1} \neq y_n$ for $n < m$ by (14) $^\bullet$. \square

We have the following analogue of Claim 4.7:

CLAIM 4.21. *Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models \text{grid}^*$. For all $m < \omega$ and all $y \in W_{x_m}$,*

- (i) if there is z such that $z \sim x_m$, $y \in W_z$ and $\mathfrak{M}, \langle z, y \rangle \models N$, then $m > 0$ and $y = y_{m-1}$,
- (ii) if there is z such that $z \sim x_m$, $y \in W_z$ and $\mathfrak{M}, \langle z, y \rangle \models \Diamond_0 N$, then $m > 1$ and $y = y_n$ for some $0 < n < m - 1$.

Proof. For (i): Take some z such that $z \sim x_m$, $y \in W_z$ and $\mathfrak{M}, \langle z, y \rangle \models N$. If $m = 0$, then $\mathfrak{M}, \langle z, y \rangle \models \blacksquare_0 \perp$ by (43) and (50), and so $\mathfrak{M}, \langle z, y \rangle \models \neg N$ by (13) $^\bullet$. So we may assume that $m > 0$. Then by (43) and Claim 4.20(ii), we have $z R^{\mathfrak{M}} x_{m-1}$, and so $y_{m-1} \in W_z$. Now (i) follows from Claim 4.20(iv) and (12) $^\bullet$.

For (ii): Take some z such that $z \sim x_m$, $y \in W_z$ and $\mathfrak{M}, \langle z, y \rangle \models \Diamond_0 N$. Then by (43), there is u such that $x_m R^{\mathfrak{M}} u$, $y \in W_u$ and $\mathfrak{M}, \langle u, y \rangle \models N$. By Claim 4.20(ii), $u \sim x_n$ for some $n < m$, and so by (i), $y = y_{n-1}$ as required. \square

Given a counter machine M , we intend to encode its runs going backward along the diagonal staircase of intervals, using again a propositional variable C_i for each $i < N$ to represent the changing content of each counter. To this end, recall the formula counter^{bw} defined in (15), and consider

$$\begin{aligned} \text{counter}^{bw\bullet} &:= \Box_1^+ \blacksquare_0 \bigwedge_{i < N} (C_i \rightarrow (N \vee \text{All} C_i^\bullet)), \quad \text{where} \\ \text{All} C_i^\bullet &:= \Diamond_0 N \wedge \blacksquare_0 (N \vee \Diamond_0 N \rightarrow C_i). \end{aligned}$$

Then we have the following analogue of Claim 4.8:

CLAIM 4.22. Suppose $\mathfrak{M}, \langle r_0, r_1 \rangle \models \text{grid}^* \wedge \text{counter}^{bw\bullet}$. Then for all $m < \omega$, $i < N$,

$$|\{y \in W_{x_{m+1}} : \mathfrak{M}, \langle I(x_{m+1}), y \rangle \models \text{AllC}_i^\bullet\}| = |\{y \in W_{x_m} : \mathfrak{M}, \langle I(x_m), y \rangle \models C_i\}|.$$

Proof. As $\blacklozenge_0 \mathbf{N}$ is a conjunct of AllC_i^\bullet , by Claim 4.21 and $\text{counter}^{bw\bullet}$, we have

$$\begin{aligned} |\{y \in W_{x_{m+1}} : \mathfrak{M}, \langle I(x_{m+1}), y \rangle \models \text{AllC}_i^\bullet\}| &= |\{n : n < m \text{ and } \mathfrak{M}, \langle I(x_{m+1}), y_n \rangle \models \text{AllC}_i^\bullet\}|, \\ |\{y \in W_{x_m} : \mathfrak{M}, \langle I(x_m), y \rangle \models C_i\}| &= |\{n : n < m \text{ and } \mathfrak{M}, \langle I(x_m), y_n \rangle \models C_i\}|. \end{aligned}$$

So it is enough to show that the two sets on the right hand sides are equal. Suppose first that $n < m$ is such that $\mathfrak{M}, \langle I(x_{m+1}), y_n \rangle \models \text{AllC}_i^\bullet$, and so

$$\mathfrak{M}, \langle x_{m+1}, y_n \rangle \models \blacksquare_0(\mathbf{N} \vee \blacklozenge_0 \mathbf{N} \rightarrow C_i).$$

Thus, in order to prove that $\mathfrak{M}, \langle I(x_m), y_n \rangle \models C_i$, it is enough to show that for all z such that $z \sim x_m$ and $y_n \in W_z$, we have

$$x_{m+1} R^{\mathfrak{M}} z \text{ and } \mathfrak{M}, \langle z, y_n \rangle \models \mathbf{N} \vee \blacklozenge_0 \mathbf{N}. \quad (51)$$

To this end, we have $x_{m+1} R^{\mathfrak{M}} x_m$ by Claim 4.20(ii), and so $x_{m+1} R^{\mathfrak{M}} z$ follows by (42). If $n = m - 1$ then $\mathfrak{M}, \langle z, y_n \rangle \models \mathbf{N}$ by Claim 4.20(iv). If $n < m - 1$ then $x_m R^{\mathfrak{M}} x_{n+1}$ by Claim 4.20(ii) and the transitivity of $R^{\mathfrak{M}}$, and so $z R^{\mathfrak{M}} x_{n+1}$ by (42). As $\mathfrak{M}, \langle x_{n+1}, y_n \rangle \models \mathbf{N}$ by Claim 4.20(iv), we obtain $\mathfrak{M}, \langle z, y_n \rangle \models \blacklozenge_0 \mathbf{N}$, as required in (51).

Conversely, suppose that $\mathfrak{M}, \langle I(x_m), y_n \rangle \models C_i$ for some $n < m$. As $n < m$, by Claims 4.20(ii),(iv) and (42), we have $\mathfrak{M}, \langle I(x_{m+1}), y_n \rangle \models \blacklozenge_0 \mathbf{N}$. In order to prove $\mathfrak{M}, \langle I(x_{m+1}), y_n \rangle \models \text{AllC}_i^\bullet$, it remains to show that

$$\mathfrak{M}, \langle I(x_{m+1}), y_n \rangle \models \blacksquare_0(\mathbf{N} \vee \blacklozenge_0 \mathbf{N} \rightarrow C_i). \quad (52)$$

To this end, let u, z be such that $u \sim x_{m+1}$, $u R^{\mathfrak{M}} z$, $y_n \in W_z$ and $\mathfrak{M}, \langle z, y_n \rangle \models \mathbf{N} \vee \blacklozenge_0 \mathbf{N}$. By (42), we have $x_{m+1} R^{\mathfrak{M}} z$, and so by (41) and Claim 4.20(ii), either $z \sim x_m$ or $x_m R^{\mathfrak{M}} z$. In the former case, $\mathfrak{M}, \langle z, y_n \rangle \models C_i$ by assumption. If $x_m R^{\mathfrak{M}} z$ then $\mathfrak{M}, \langle x_m, y_n \rangle \models \neg \mathbf{N}$ by (13) $^\bullet$, and so $\mathfrak{M}, \langle x_m, y_n \rangle \models \text{AllC}_i^\bullet$ by $\text{counter}^{bw\bullet}$. Thus, $\mathfrak{M}, \langle x_m, y_n \rangle \models \blacksquare_0(\mathbf{N} \vee \blacklozenge_0 \mathbf{N} \rightarrow C_i)$, and so $\mathfrak{M}, \langle z, y_n \rangle \models C_i$ follows in this case as well, proving (52). \square

Now recall the formulas Fix_i^{bw} , Inc_i^{bw} and Dec_i^{bw} from (16)–(18), simulating the possible changes in the counters stepping backward, and ensuring that each ‘vertical coordinate’ is used only once in the counting. Observe that $\square_1^+ \square_0^+(C_i \rightarrow \blacksquare_0 \neg C_i)$ (conjunct (44) of Interval_{C_i}) and $\text{counter}^{bw\bullet}$ cannot hold simultaneously, so we cannot use the formula Interval_{C_i} for forcing C_i to behave uniformly in intervals. However, as each vertical coordinate is used at most once in the counting, we can force that the *changes* happen uniformly in the intervals (even when the counter is decremented). To this end, for each $i < N$ we introduce a fresh propositional variable C_i^- , and then postulate

$$\bigwedge_{i < N} \left(\text{Interval}_{C_i^-} \wedge \square_1^+ \square_0^+(C_i^- \leftrightarrow (\neg C_i \wedge \text{AllC}_i^\bullet)) \right). \quad (53)$$

Now we have the following analogue of Claim 4.9:

CLAIM 4.23. Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models \text{grid}^* \wedge \text{counter}^{bw\bullet} \wedge (53)$ and, for all $m < \omega$, $i < n$, let $c_i(m) := |\{y \in W_{x_m} : \mathfrak{M}, \langle I(x_m), y \rangle \models C_i\}|$. Then

$$c_i(m+1) = \begin{cases} c_i(m), & \text{if } \mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models \text{Fix}_i^{bw\bullet}, \\ c_i(m) + 1, & \text{if } \mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models \text{Inc}_i^{bw\bullet}, \\ c_i(m) - 1, & \text{if } \mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models \text{Dec}_i^{bw\bullet}. \end{cases}$$

Proof. We show only the hardest case, when $\mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models \text{Dec}_i^{bw\bullet}$. The other cases are similar and left to the reader. As $\mathfrak{M}, \langle x_{m+1}, y_{m+1} \rangle \models \Diamond_1^{-1}(\neg C_i \wedge \text{AllC}_i^\bullet)$, there is an $y^* \in W_{x_{m+1}}$ such that

$$\mathfrak{M}, \langle x_{m+1}, y^* \rangle \models \neg C_i \wedge \text{AllC}_i^\bullet, \quad (54)$$

$$\mathfrak{M}, \langle x_{m+1}, y \rangle \not\models \neg C_i \wedge \text{AllC}_i^\bullet, \text{ for all } y \neq y^*, y \in W_{x_{m+1}}. \quad (55)$$

We claim that

$$\{y \in W_{x_{m+1}} : \mathfrak{M}, \langle I(x_{m+1}), y \rangle \models C_i\} \cup \{y^*\} = \{y \in W_{x_{m+1}} : \mathfrak{M}, \langle I(x_{m+1}), y \rangle \models \text{AllC}_i^\bullet\}. \quad (56)$$

Indeed, in order to show the \subseteq direction, suppose first that $\mathfrak{M}, \langle I(x_{m+1}), y \rangle \models C_i$ for some $y \in W_{x_{m+1}}$. Then by the first conjunct of $\text{Dec}_i^{bw\bullet}$, we have $\mathfrak{M}, \langle I(x_{m+1}), y \rangle \models \text{AllC}_i^\bullet$. Further, we have $\mathfrak{M}, \langle I(x_{m+1}), y^* \rangle \models \text{AllC}_i^\bullet$ by (54), (53) and Claim 4.19. For \supseteq , suppose that $\mathfrak{M}, \langle I(x_{m+1}), y \rangle \models \text{AllC}_i^\bullet$ for some $y \in W_{x_{m+1}}, y \neq y^*$. Then by (55), (53) and Claim 4.19, we have $\mathfrak{M}, \langle I(x_{m+1}), y \rangle \models \neg C_i^-$, and so $\mathfrak{M}, \langle I(x_{m+1}), y \rangle \models C_i$, proving (56).

Now $c_i(m+1) + 1 = c_i(m)$ follows from (56) and Claim 4.22. \square

Given a counter machine M , recall the formula φ_M^{bw} defined in the proof of Theorem 4.1 (as the conjunction of (15) and (19)–(21)). Let φ_M^* be the conjunction of $\varphi_M^{bw\bullet}$, (53) and Interval_P , for $P \in \{S_q, I_\iota\}_{q \in Q, \iota \in \text{Op}_C}$. Then we have the following analogue of Lemma 4.10:

Lemma 4.24. *Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models \text{grid}^* \wedge \varphi_M^*$, and for all $m < \omega, i < N$, let*

$$q_m := q, \text{ if } \mathfrak{M}, \langle I(x_m), y_m \rangle \models S_q, \quad c_i(m) := |\{y \in W_{x_m} : \mathfrak{M}, \langle I(x_m), y \rangle \models C_i\}|.$$

Then $\langle \langle q_m, c(m) \rangle : m < \omega \rangle$ is a well-defined infinite run of M starting with all-0 counters.

Proof. The sequence $\langle q_m : m < \omega \rangle$ is well-defined by Claims 4.20(iii), 4.19 and (19) $^\bullet$. We show by induction on m that for all $m < \omega$, $\langle \langle q_0, c(0) \rangle, \dots, \langle q_m, c(m) \rangle \rangle$ is a run of M starting with all-0 counters. Indeed, $c_i(0) = 0$ for $i < N$ by **counter** $^{bw\bullet}$ and Claim 4.21. Now suppose the statement holds for some $m < \omega$. By Claim 4.20, $\mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models S \wedge \Diamond_1(N \wedge \Diamond_0 S_{q_m})$. So by (20) $^\bullet$, there is $\langle \iota, q_{m+1} \rangle \in I_{q_m}$ such that $\mathfrak{M}, \langle x_{m+1}, y_{m+1} \rangle \models I_\iota \wedge S_{q_{m+1}}$, and so $\mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models I_\iota$ by Interval_I and Claim 4.19. Thus, $\mathfrak{M}, \langle I(x_{m+1}), y_{m+1} \rangle \models \text{Do}_\iota^{bw\bullet}$ by (21) $^\bullet$. It follows from Claim 4.23 that $\sigma_m \rightarrow^\iota \sigma_{m+1}$ as required. \square

For the other direction, suppose that M has an infinite run starting with all-0 counters. Let $\mathfrak{F} = \langle W, R \rangle$ be a weak order in \mathcal{C} containing an $\langle \omega + 1, > \rangle$ -type chain $x_\omega R \dots R x_m R \dots R x_0$. For every $m < \omega$, we let

$$[x_{m+1}, x_m] := (\{w \in W : x_{m+1} R w R x_m\} \cup \{x_{m+1}\}) - \{w : w = x_m \text{ or } x_m R w\}.$$

Take the model $\mathfrak{M}^\infty = \langle \langle \omega + 1, > \rangle \times \langle \omega, \neq \rangle, \mu \rangle$ defined in (22)–(27). We define a model $\mathfrak{N}^\infty = \langle \mathfrak{F} \times \langle \omega, \neq \rangle, \nu \rangle$ as follows. We let

$$\nu(\text{Tick}) := \{\langle w, n \rangle : w \in [x_{m+1}, x_m], m, n < \omega, m \text{ is odd}\},$$

for all $P \in \{N, S, S_q, I_\iota, C_i\}_{q \in Q, \iota \in \text{Op}_C, i < N}$,

$$\nu(P) := \{\langle w, n \rangle : w \in [x_{m+1}, x_m], \langle m, n \rangle \in \mu(P) \text{ for some } m < \omega\},$$

for all $P \in \{N, S, S_q, I_\iota\}_{q \in Q, \iota \in \text{Op}_C}$,

$$\nu(P') := \{\langle w, n \rangle : w \in [x_m, x_{m-1}], \langle m, n \rangle \in \mu(P) \text{ for some } m > 0\},$$

and for all $i < N$,

$$\nu(C_i^-) := \{\langle w, n \rangle : w \in [x_{m+1}, x_m], \langle m, n \rangle \notin \mu(C_i), \langle m-1, n \rangle \in \mu(C_i) \text{ for some } m > 0\},$$

$$\nu(C_i'^-) := \{\langle w, n \rangle : w \in [x_{m+1}, x_m], \langle m, n \rangle \in \mu(C_i), \langle m+1, n \rangle \notin \mu(C_i) \text{ for some } m < \omega\}.$$

It is not hard to check that $\mathfrak{N}^\infty, \langle x_\omega, 0 \rangle \models \text{grid}^* \wedge \varphi_M^*$. So by Lemma 4.24, CM non-termination is reducible to $\mathcal{C} \times {}^d\mathcal{C}_{\text{diff}}$ -satisfiability. This completes the proof of Theorem 4.16.

5 Expanding 2-frames

In this section we show that satisfiability over classes of expanding 2-frames can be genuinely simpler than satisfiability over the corresponding product frame classes, but it is still quite complex.

5.1 Lower bounds

Theorem 5.1. $\{\langle \omega, < \rangle\} \times {}^e\mathcal{C}_{diff}$ -satisfiability is undecidable.

Corollary 5.2. *FOLTL[≠]-satisfiability is undecidable in expanding domain models over $\langle \omega, < \rangle$.*

Theorem 5.3. $\mathcal{C}_{lin}^{fin} \times^e \mathcal{C}_{diff}$ -satisfiability is Ackermann-hard.

Corollary 5.4. *FOLTE[≠]-satisfiability is Ackermann-hard in expanding domain models over the class of all finite linear orders.*

We prove Theorem 5.1 by reducing the ‘LCM ω -reachability’ problem to $\{\langle \omega, < \rangle\} \times {}^e\mathcal{C}_{diff}$ -satisfiability. The idea of our reduction is similar to the one used in [27] for a more expressive formalism. It is sketched in Fig. 4: First, we generate an infinite diagonal staircase going forward. Then, still going forward, we place longer and longer finite runs one after the other. However, each individual run proceeds backward. Also, we can force only lossy runs this way. When going backward horizontally in expanding 2-frames, the vertical columns might become smaller and smaller, so some of the points carrying the information on the content of the counters might disappear as the runs progress.

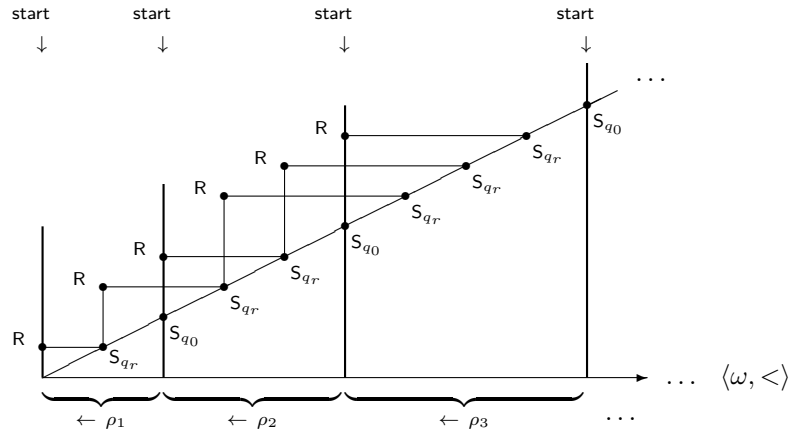


Figure 4: Representing longer and longer n -recurrent lossy runs ρ_n in 2-frames expanding over $\langle \omega, < \rangle$.

To this end, let $\mathfrak{H}_{\langle\omega, <\rangle, \overline{\mathfrak{G}}}$ be an expanding 2-frame for some difference frames $\mathfrak{G}_n = \langle W_n, \neq \rangle$, $n < \omega$, and let \mathfrak{M} be a model based on $\mathfrak{H}_{\langle\omega, <\rangle, \overline{\mathfrak{G}}}$. First, we generate an infinite diagonal staircase forward in \mathfrak{M} , similarly how we did in the proof of Theorem 3.1. However, this time we use the vertical counting capabilities to force the *uniqueness* of this staircase. To this end, let `grid_unique` be the conjunction of (3)–(5) and

$$\square_0^+ \square_1 (\mathbf{N} \rightarrow \square_1 \neg \mathbf{N}). \quad (57)$$

The following ‘expanding generalisation’ of Claim 3.5 can be proved by a straightforward induction on m :

CLAIM 5.5. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique}$. Then there exists a sequence $\langle y_m : m < \omega \rangle$ such that for all $m < \omega$,*

- (i) $y_0 = r$ and if $m > 0$ then $y_m \in W_{m-1}$,
- (ii) for all $n < m$, $y_m \neq y_n$,
- (iii) $\mathfrak{M}, \langle m, y_m \rangle \models S$,
- (iv) for all $w \in W_m$, $\mathfrak{M}, \langle m, w \rangle \models N$ iff $w = y_{m+1}$.

Given a counter machine M , we will encode lossy runs that start with all-0 counters by going backward along the created diagonal staircase. We will adjust the tools developed in the proof of Theorem 4.1 in order to handle lossiness, and also to force not just one run, but several (finite) runs, placed one after the other. To this end, we introduce a fresh propositional variable **start**, intended to mark the start of each run (see Fig. 4), and for each $i < N$ we let

$$\text{TillStartAllC}_i := \Diamond_0 N \wedge \Box_0 (N \vee \Diamond_0 N \rightarrow (\neg \text{start} \wedge C_i)).$$

Then we have the following lossy analogue of Claims 4.8 and 4.22:

CLAIM 5.6. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique}$. Then for all $m < \omega$, $i < N$,*

$$\{w \in W_m : \mathfrak{M}, \langle m, w \rangle \models \text{TillStartAllC}_i\} \subseteq \{w \in W_{m+1} : \mathfrak{M}, \langle m+1, w \rangle \models C_i\}.$$

Proof. Suppose that $\mathfrak{M}, \langle m, w \rangle \models \text{TillStartAllC}_i$. Then $\mathfrak{M}, \langle m, w \rangle \models \Diamond_0 N$ and so by Claim 5.5(iv), $w = y_n$ for some $n > m+1$, and we have $\mathfrak{M}, \langle n-1, w \rangle \models N$. Thus, $\mathfrak{M}, \langle m+1, w \rangle \models N \vee \Diamond_0 N$. As $\mathfrak{M}, \langle m, w \rangle \models \Box_0 (N \vee \Diamond_0 N \rightarrow C_i)$, we obtain $\mathfrak{M}, \langle m+1, w \rangle \models C_i$ as required. \square

Now, for each $i < N$, we can simulate the possible lossy changes in the counters by the following formulas:

$$\begin{aligned} \text{Fix}_i^{\text{lossy}} &:= \Box_1^+(C_i \rightarrow \text{TillStartAllC}_i), \\ \text{Inc}_i^{\text{lossy}} &:= \Box_1^+(C_i \rightarrow (N \vee \text{TillStartAllC}_i)), \\ \text{Dec}_i^{\text{lossy}} &:= \Box_1^+(C_i \rightarrow \text{TillStartAllC}_i) \wedge \Box_1^+(\neg C_i \wedge \text{TillStartAllC}_i). \end{aligned}$$

The following lossy analogue of Claims 4.9 and 4.23 is a straightforward consequence of Claims 5.5(iv) and 5.6. Note that the vertical uniqueness of N -points is used in simulating the lossy incrementation steps properly.

CLAIM 5.7. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique}$. For all $i < N$, $m < \omega$, let $c_i(m) := |\{w \in W_m : \mathfrak{M}, \langle m, w \rangle \models C_i\}|$. Then for all $m < \omega$,*

$$c_i(m) \leq \begin{cases} c_i(m+1), & \text{if } \mathfrak{M}, \langle m, y_m \rangle \models \text{Fix}_i^{\text{lossy}}, \\ c_i(m+1) + 1, & \text{if } \mathfrak{M}, \langle m, y_m \rangle \models \text{Inc}_i^{\text{lossy}}, \\ c_i(m+1) - 1, & \text{if } \mathfrak{M}, \langle m, y_m \rangle \models \text{Dec}_i^{\text{lossy}}. \end{cases}$$

Next, we encode the various counter machine instructions for lossy steps, acting backward. For each $\iota \in \text{Op}_C$, we define the formula $\text{Do}_\iota^{\text{lossy}}$ by taking

$$\text{Do}_\iota^{\text{lossy}} := \begin{cases} \text{Inc}_i^{\text{lossy}} \wedge \bigwedge_{i \neq j < N} \text{Fix}_j^{\text{lossy}}, & \text{if } \iota = c_i^{++}, \\ \text{Dec}_i^{\text{lossy}} \wedge \bigwedge_{i \neq j < N} \text{Fix}_j^{\text{lossy}}, & \text{if } \iota = c_i^{--}, \\ \Box_1^+ \neg C_i \wedge \bigwedge_{j < N} \text{Fix}_j^{\text{lossy}}, & \text{if } \iota = c_i^{??}. \end{cases}$$

Finally, given a counter machine M , we encode lossy runs that start with all-0 counters at **start**-marks, and go backward until the next **start**-mark. We define φ_M^{lossy} to be the conjunction of the following formulas:

$$\Box_0^+ \Box_1^+ (\text{start} \rightarrow \Box_1 \text{start}), \quad (58)$$

$$\Box_0^+ \Box_1^+ (S \leftrightarrow \bigvee_{q \in Q-H} (S_q \wedge \bigwedge_{q \neq q' \in Q} \neg S_{q'})), \quad (59)$$

$$\Box_0^+ \Box_1^+ (S \wedge \text{start} \rightarrow (S_{q_0} \wedge \bigwedge_{i < N} \Box_1^+ \neg C_i)), \quad (60)$$

$$\Box_0 \Box_1 \bigwedge_{q \in Q-H} [(S \wedge \neg \text{start} \wedge \Diamond_1 (N \wedge \Diamond_0 S_q)) \rightarrow \bigvee_{\langle \iota, q' \rangle \in I_q} (\text{Do}_\iota^{lossy} \wedge S_{q'})]. \quad (61)$$

Then we have the following lossy analogue of Lemmas 4.10 and 4.24:

CLAIM 5.8. *Suppose $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique} \wedge \varphi_M^{lossy}$, and for all $m < \omega$, $i < N$, let*

$$s_m := q, \text{ if } \mathfrak{M}, \langle m, y_m \rangle \models S_q, \quad c_i(m) := |\{w \in W_m : \mathfrak{M}, \langle m, w \rangle \models C_i\}|, \quad \sigma_m := \langle s_m, \mathbf{c}(m) \rangle.$$

Then $\langle \sigma_a, \sigma_{a-1}, \dots, \sigma_b \rangle$ is a well-defined lossy run of M starting with $\langle q_0, \mathbf{0} \rangle$, whenever $b < a < \omega$ is such that $\mathfrak{M}, \langle a, r \rangle \models \text{start}$, and $\mathfrak{M}, \langle n, r \rangle \models \neg \text{start}$, for every n with $b \leq n < a$.

Proof. The sequence $\langle s_a, s_{a-1}, \dots, s_b \rangle$ is well-defined by Claim 5.5(iii) and (59). We show by induction on m that for all $m \leq a - b$, $\langle \sigma_a, \sigma_{a-1}, \dots, \sigma_{a-m} \rangle$ is a lossy run of M starting with $\langle q_0, \mathbf{0} \rangle$. Indeed, $\mathfrak{M}, \langle a, y_a \rangle \models \text{start}$ by (58), and so $s_a = q_0$ and $c_i(a) = 0$ for $i < N$ by Claim 5.5(iii) and (60). Now suppose the statement holds for some $m < a - b$. As $\mathfrak{M}, \langle a - m - 1, y_{a-m-1} \rangle \models \neg \text{start}$ by (58), we have

$$\mathfrak{M}, \langle a - m - 1, y_{a-m-1} \rangle \models S \wedge \neg \text{start} \wedge \Diamond_1 (N \wedge \Diamond_0 S_{s_{a-m}})$$

by Claim 5.5. By (59) we have $s_{a-m} \in Q - H$, and so by (61) there is $\langle \iota, s_{a-m-1} \rangle \in I_{s_{a-m}}$ such that $\mathfrak{M}, \langle a - m - 1, y_{a-m-1} \rangle \models \text{Do}_\iota^{lossy}$. It follows from Claims 5.6 and 5.7 that $\sigma_{a-m} \xrightarrow{\iota}_{lossy} \sigma_{a-m-1}$ as required. \square

It remains to force that the n th run visits q_r at least n times. To this end, we introduce two fresh propositional variables R and S^* , and define rec as the conjunction of (58) and the following formulas:

$$\text{start} \wedge \Box_0^+ \Diamond_0 \text{start}, \quad (62)$$

$$\Box_0^+ \Box_1^+ [\text{start} \rightarrow \Diamond_1^+ (R \wedge \Diamond_0 (S \wedge \neg \text{start}) \wedge \Box_0 (\Diamond_0 S \rightarrow \neg \text{start}))], \quad (63)$$

$$\Box_0^+ \Box_1^+ (R \rightarrow \Box_0 (S \rightarrow S^*)), \quad (64)$$

$$\Box_0 \Box_1^+ [S^* \rightarrow \Diamond_1 [R \wedge \Diamond_0 (\text{start} \wedge \Diamond_0 (S \wedge \neg \text{start})) \wedge \Box_0 (\text{start} \wedge \Diamond_0 S \rightarrow \Box_0 (\Diamond_0 S \rightarrow \neg \text{start}))]], \quad (65)$$

$$\Box_0^+ \Box_1^+ (S^* \rightarrow S), \quad (66)$$

$$\Box_0^+ \Box_1^+ (S \rightarrow \Box_1 \neg S), \quad (67)$$

$$\Box_0^+ \Box_1^+ (R \rightarrow \Box_0 \neg R). \quad (68)$$

CLAIM 5.9. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique} \wedge \text{rec}$. Then there is an infinite sequence $\langle k_n : n < \omega \rangle$ such that, for all $n < \omega$,*

(i) $\mathfrak{M}, \langle k_n, w \rangle \models \text{start}$ for all $w \in W_{k_n}$,

(ii) if $n > 0$ then $\mathfrak{M}, \langle k, w \rangle \models \neg \text{start}$ for all k with $k_{n-1} < k < k_n$ and $w \in W_k$, and

(iii) if $n > 0$ then $|\{k : k_{n-1} < k < k_n \text{ and } \mathfrak{M}, \langle k, y_k \rangle \models S^*\}| \geq n$.

Proof. By induction on n . To begin with, let $k_0 = 0$. Now suppose inductively that we have $\langle k_\ell : \ell < n \rangle$ as required, for some $0 < n < \omega$. Now let k_n be the smallest k with $k > k_{n-1}$ and $\mathfrak{M}, \langle k, r \rangle \models \text{start}$ (there is such by (62)). So $k_n > k_{n-1}$, and by (58)

$$\mathfrak{M}, \langle k_n, w \rangle \models \text{start} \text{ for all } w \in W_{k_n}. \quad (69)$$

As by the IH(i) we have $\mathfrak{M}, \langle k_{n-1}, r \rangle \models \text{start}$, by (63) there is $w \in W_{k_{n-1}}$ such that

$$\mathfrak{M}, \langle k_{n-1}, w \rangle \models R \wedge \Diamond_0(S \wedge \neg \text{start}) \wedge \Box_0(\Diamond_0 S \rightarrow \neg \text{start}).$$

By Claim 5.5(iii) and (67), $w = y_{i_n}$ for some $k_{n-1} < i_n < k_n$, and so $\mathfrak{M}, \langle i_n, y_{i_n} \rangle \models S^*$ follows by (64). In particular, if $n = 1$ then $\mathfrak{M}, \langle i_1, y_{i_1} \rangle \models S^*$, and so

$$|\{k : k_0 < k < k_1 \text{ and } \mathfrak{M}, \langle k, y_k \rangle \models S^*\}| \geq 1.$$

Now suppose that $n > 1$ and take some k such that $k_{n-2} < k < k_{n-1}$ and $\mathfrak{M}, \langle k, y_k \rangle \models S^*$. By (65), there is $v \in W_k$ such that

$$\mathfrak{M}, \langle k, v \rangle \models R \wedge \Diamond_0(\text{start} \wedge \Diamond_0(S \wedge \neg \text{start})) \wedge \Box_0(\text{start} \wedge \Diamond_0 S \rightarrow \Box_0(\Diamond_0 S \rightarrow \neg \text{start})). \quad (70)$$

So there is some $k' > k$ with $\mathfrak{M}, \langle k', v \rangle \models \text{start} \wedge \Diamond_0(S \wedge \neg \text{start})$, and so by the IH we have

$$\mathfrak{M}, \langle k_{n-1}, v \rangle \models \text{start} \wedge \Diamond_0(S \wedge \neg \text{start}). \quad (71)$$

Therefore, by (70) we have

$$\mathfrak{M}, \langle k_{n-1}, v \rangle \models \Box_0(\Diamond_0 S \rightarrow \neg \text{start}). \quad (72)$$

By (71), there is some $k^+ > k_{n-1}$ with $\mathfrak{M}, \langle k^+, v \rangle \models S \wedge \neg \text{start}$. Therefore, $v = y_{k^+}$ by Claim 5.5(iii) and (67), $\mathfrak{M}, \langle k^+, y_{k^+} \rangle \models S^*$ by (64), and $k^+ \neq k_n$ by (69). Moreover, we have that $k^+ < k_n$ because of the following. If $k^+ > k_n$ were the case, then $\mathfrak{M}, \langle k_n, v \rangle \models \Diamond_0 S$, and so $\mathfrak{M}, \langle k_n, v \rangle \models \neg \text{start}$ by (72), contradicting (69). Further, by (68) we obtain that $k^+ \neq i_n$, and $k^+ \neq \ell^+$ whenever $k \neq \ell$, $k_{n-1} < k, \ell < k_n$. Therefore, by (66), (67), and the IH(iii), we have

$$\begin{aligned} & |\{k : k_{n-1} < k < k_n \text{ and } \mathfrak{M}, \langle k, y_k \rangle \models S^*\}| \geq \\ & |\{k : k_{n-2} < k < k_{n-1} \text{ and } \mathfrak{M}, \langle k, y_k \rangle \models S^*\}| + 1 \geq n - 1 + 1 = n, \end{aligned}$$

as required. \square

Now the following lemma is a straightforward consequence of Claims 5.8 and 5.9:

Lemma 5.10. *Suppose $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique} \wedge \varphi_M^{\text{lossy}} \wedge \text{rec} \wedge \Box_0^+ \Box_1^+(S^* \rightarrow S_{q_r})$. Then, for every $n < \omega$, M has a lossy run starting with $\langle q_0, \mathbf{0} \rangle$ and visiting q_r at least n times.*

On the other hand, suppose that for every $0 < n < \omega$, M has a lossy run

$$\rho_n = \langle \langle q_0^n, \mathbf{0} \rangle, \dots, \langle q_{m_n-1}^n, \mathbf{c}(m_n - 1) \rangle \rangle$$

such that $q_0^n = q_0$ and ρ_n visits q_r at least n times. Let $M_0 := 0$ and for each $0 < n < \omega$, let $M_n := \sum_{i=1}^n m_i$, and let $i_1^n, \dots, i_n^n < m_n$ be such that $|\{i_1^n, \dots, i_n^n\}| = n$ and $q_i = q_r$ for every $i \in \{i_1^n, \dots, i_n^n\}$. We define a model $\mathfrak{M}^\infty = \langle \langle \omega, < \rangle \times \langle \omega, \neq \rangle, \alpha \rangle$ as follows (cf. Fig. 4): For all $q \in Q$, we let

$$\begin{aligned} \alpha(S_q) &:= \{ \langle n, n \rangle : M_k \leq n < M_{k+1} \text{ and } q_{n-M_k}^{k+1} = q, \text{ for some } k < \omega \}, \\ \alpha(S) &:= \{ \langle n, n \rangle : n < \omega \}, \\ \alpha(\mathbf{N}) &:= \{ \langle n, n+1 \rangle : n < \omega \}, \\ \alpha(\text{start}) &:= \{ \langle n, m \rangle : n = M_k \text{ for some } k < \omega, \text{ and } m < \omega \}. \end{aligned}$$

Further, for any finite subset $X = \{n_1, \dots, n_\ell\}$ of ω with $n_1 < \dots < n_\ell$ and any $k \leq |X|$, we let $\min_k(X) := \{n_1, \dots, n_k\}$. Now for all $i < N$, $0 < n < \omega$ and $k < m_n$, we define the sets $\alpha_k^n(C_i)$ by induction on k : We let $\alpha_0^n(C_i) := \emptyset$, and for all $k < m_n - 1$,

$$\alpha_{k+1}^n(C_i) := \begin{cases} \alpha_k^n(C_i) \cup \{M_n - k\}, & \text{if } c_i^n(k+1) = c_i^n(k) + 1, \\ \alpha_k^n(C_i) - \min_\ell(\alpha_k^n(C_i)), & \text{if } |c_i^n(k) - c_i^n(k+1)| = \ell. \end{cases}$$

Then, for each $i < N$, we let

$$\alpha(C_i) := \{\langle k, m \rangle : M_{n-1} \leq k < M_n, m \in \alpha_{M_n-k-1}^n(C_i) \text{ for some } 0 < n < \omega\}.$$

Also, we define the sequence $\langle r_n : n < \omega \rangle$ inductively as follows. Let $r_0 := i_1^1$ and let

$$r_{n+1} := \begin{cases} M_k + i_1^{k+1}, & \text{if } r_n = M_{k-1} + i_k^k \text{ for some } k > 0, \\ M_{k-1} + i_{\ell+1}^k, & \text{if } r_n = M_{k-1} + i_\ell^k \text{ for some } k > 0, \ell < k. \end{cases}$$

Then let

$$\begin{aligned} \alpha(S^*) &:= \{\langle r_n, r_n \rangle : n < \omega\}, \\ \alpha(R) &:= \{\langle n, r_n \rangle : n < \omega\}. \end{aligned}$$

It is not hard to check that $\mathfrak{N}^\infty, \langle 0, 0 \rangle \models \text{grid_unique} \wedge \varphi_M^{\text{lossy}} \wedge \text{rec} \wedge \Box_0^+ \Box_1^+(S^* \rightarrow S_{q_r})$, and so by Lemma 5.10 LCM ω -reachability can be reduced to $\{\langle \omega, < \rangle\} \times^e \mathcal{C}_{\text{diff}}$ -satisfiability. This completes the proof of Theorem 5.1.

Next, we prove Theorem 5.3 by reducing the ‘LCM-reachability’ problem to $\mathcal{C}_{\text{lin}}^{\text{fin}} \times^e \mathcal{C}_{\text{diff}}$ -satisfiability. We will use the finitary versions of some of the formulas used in the previous proof. Let $\mathfrak{H}_{\langle T, < \rangle, \overline{\mathfrak{G}}}$ be an expanding 2-frame for some finite linear order $\langle T, < \rangle$ and for some difference frames $\mathfrak{G}_n = \langle W_n, \neq \rangle$, $n \in T$, and let \mathfrak{M} be a model based on $\mathfrak{H}_{\langle T, < \rangle, \overline{\mathfrak{G}}}$. We may assume that $T = |T| < \omega$. We consider a version of the formula grid_{fin} defined in the proof of Theorem 3.2. Let $\text{grid_unique}_{\text{fin}}$ be the conjunction of (3), (4), (9) and (57). The following finitary version of Claim 5.5 can be proved by a straightforward induction on m :

CLAIM 5.11. *Suppose $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique}_{\text{fin}}$. Then there exist some $0 < E \leq T$ and a sequence $\langle y_m : m \leq E \rangle$ of points such that for all $m \leq E$,*

- (i) $y_0 = r$ and if $m > 0$ then $y_m \in W_{m-1}$,
- (ii) for all $n < m$, $y_m \neq y_n$,
- (iii) if $m < E$ then $\mathfrak{M}, \langle m, y_m \rangle \models S$,
- (iv) if $m < E$ then for all $w \in W_m$, $\mathfrak{M}, \langle m, w \rangle \models N$ iff $w = y_{m+1}$,
- (v) $\mathfrak{M}, \langle E-1, y_E \rangle \models \text{end}$, and if $m < E-1$ then $\mathfrak{M}, \langle m, y_{m+1} \rangle \models \neg \text{end}$.

The following lemma is a straightforward consequence of Claims 5.8 and 5.11:

Lemma 5.12. *Suppose that $\mathfrak{M}, \langle 0, r \rangle \models \text{grid_unique}_{\text{fin}} \wedge \varphi_M^{\text{lossy}} \wedge S_{q_r} \wedge \Box_0^+ \Box_1^+(\text{end} \leftrightarrow \text{start})$. For all $m < E$ and $i < N$, let*

$$s_m := q, \text{ if } \mathfrak{M}, \langle m, y_m \rangle \models S_q, \quad c_i(m) := |\{w \in W : \mathfrak{M}, \langle m, w \rangle \models C_i\}|, \quad \sigma_m = \langle s_m, \mathbf{c}(m) \rangle.$$

Then $\langle \sigma_{E-1}, \sigma_{E-2}, \dots, \sigma_0 \rangle$ is a well-defined lossy run of M starting with $\langle q_0, \mathbf{0} \rangle$ and reaching q_r .

On the other hand, if M has a run $\langle \langle q_m, \mathbf{c}(m) \rangle : m < T \rangle$ for some $T < \omega$ such that it starts with all-0 counters and $q_{T-1} = q_r$, then it is not hard to define a model based on $\langle T, < \rangle \times \langle T+1, \neq \rangle$ satisfying $\text{grid_unique}_{\text{fin}} \wedge \varphi_M^{\text{lossy}} \wedge S_{q_r} \wedge \Box_0^+ \Box_1^+(\text{end} \leftrightarrow \text{start})$ (cf. how the finite runs in the model \mathfrak{N}^∞ are defined in the proof of Theorem 5.1). So by Lemma 5.12 the proof of Theorem 5.3 is completed.

5.2 Upper bounds

To begin with, as a consequence of Theorems 2.1 and Props. 2.2, 2.3 we obtain:

Corollary 5.13. *FOLTL[≠]-satisfiability is co-r.e. in expanding domain models over the class of all linear orders.*

Unlike in the constant domain case, in the expanding domain case the same holds for $\langle \omega, < \rangle$ as timeline:

Theorem 5.14. *$\{\langle \omega, < \rangle\} \times {}^e\mathcal{C}_{diff}$ -satisfiability is co-r.e.*

Corollary 5.15. *FOLTL[≠]-satisfiability is co-r.e. in expanding domain models over $\langle \omega, < \rangle$.*

Theorem 5.16. *$\mathcal{C}_{lin}^{fin} \times {}^e\mathcal{C}_{diff}$ -satisfiability is decidable.*

Corollary 5.17. *FOLTL[≠]-satisfiability is decidable in expanding domain models over the class of all finite linear orders.*

In order to prove both Theorems 5.14 and 5.16, we begin with showing that there is a reduction from \mathcal{C}_{diff} -satisfiability to \mathcal{C}_{lin} -satisfiability that can be ‘lifted to the 2D level’. As we will use this reduction to obtain upper bounds on satisfiability in expanding 2-frames, we formulate it in this setting only. To this end, fix some bimodal formula ϕ . For every $\psi \in \text{sub } \phi$, we introduce a fresh propositional variable P_ψ not occurring in ϕ , and define inductively a translation ψ^\dagger by taking

$$\begin{aligned} P^\dagger &:= P, \text{ for each propositional variable } P \in \text{sub } \phi, \\ (\neg\psi)^\dagger &:= \neg\psi^\dagger, \\ (\psi_1 \wedge \psi_2)^\dagger &:= \psi_1^\dagger \wedge \psi_2^\dagger, \\ (\Diamond_0\psi)^\dagger &:= \Diamond_0\psi^\dagger, \\ (\Diamond_1\psi)^\dagger &:= P_\psi \vee \Diamond_1\psi^\dagger. \end{aligned}$$

Further, we let

$$\chi_\phi := \Box_0^+ \bigwedge_{\psi \in \text{sub } \phi} \neg P_\psi \wedge \Box_1^+(\psi^\dagger \rightarrow \Box_1 P_\psi) \wedge (\Diamond_1 P_\psi \rightarrow \Diamond_1^+(\neg P_\psi \wedge \psi^\dagger)).$$

CLAIM 5.18. *For any formula ϕ , and any class \mathcal{C} of transitive frames,*

- *ϕ is $\mathcal{C} \times {}^e\mathcal{C}_{diff}$ -satisfiable iff $\chi_\phi \wedge \phi^\dagger$ is $\mathcal{C} \times {}^e\mathcal{C}_{lin}$ -satisfiable.*
- *ϕ is $\mathcal{C} \times {}^e\mathcal{C}_{diff}^{fin}$ -satisfiable iff $\chi_\phi \wedge \phi^\dagger$ is $\mathcal{C} \times {}^e\mathcal{C}_{lin}^{fin}$ -satisfiable.*

Proof. \Rightarrow : Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models \phi$ in some model $\mathfrak{M} = \langle \mathfrak{H}_{\mathfrak{F}, \overline{G}}, \mu \rangle$ based on an expanding 2-frame $\mathfrak{H}_{\mathfrak{F}, \overline{G}}$ where $\mathfrak{F} = \langle W, R \rangle$ is transitive and for every $x \in W$, $\mathfrak{G}_x = \langle W_x, \neq \rangle$. Then $W_x \subseteq W_y$ whenever xRy , $x, y \in W$. Also, we may assume that r_0 is a root in \mathfrak{F} , and so $r_1 \in W_x$ for all $x \in W$. So for every $x \in W$ we may take a well-order $<_x$ on W_x with least element r_1 and such that $<_x \subseteq <_y$ whenever xRy . Let $\Sigma'_x = \langle W_x, <_x \rangle$, for $x \in W$. Then clearly $\mathfrak{H}_{\mathfrak{F}, \overline{G}'} \in \mathcal{C} \times {}^e\mathcal{C}_{lin}$. We define a model $\mathfrak{M}' = \langle \mathfrak{H}_{\mathfrak{F}, \overline{G}'}, \mu' \rangle$ by taking

$$\begin{aligned} \mu'(P) &:= \mu(P), \text{ for } P \in \text{sub } \phi, \\ \mu'(P_\psi) &:= \{ \langle x, w \rangle : x \in W \text{ and } \mathfrak{M}, \langle x, u \rangle \models \psi \text{ for some } u \in W_x \text{ with } u <_x w \}. \end{aligned}$$

First, we show by induction on ψ that for all $\psi \in \text{sub } \phi$, $x \in W$, $u \in W_x$,

$$\mathfrak{M}, \langle x, u \rangle \models \psi \quad \text{iff} \quad \mathfrak{M}', \langle x, u \rangle \models \psi^\dagger. \quad (73)$$

Indeed, the only non-straightforward case is that of \Diamond_1 . So suppose first that $\mathfrak{M}, \langle x, u \rangle \models \Diamond_1\psi$. Then there is $v \in W_x$, $v \neq u$ with $\mathfrak{M}, \langle x, v \rangle \models \psi$. If $u <_x v$ then $\mathfrak{M}', \langle x, u \rangle \models \Diamond_1\psi^\dagger$ by the IH. If

$v <_x u$, then $\mathfrak{M}', \langle x, u \rangle \models P_\psi$ by the definition of \mathfrak{M}' . So in both cases we have $\mathfrak{M}', \langle x, u \rangle \models (\Diamond_1 \psi)^\dagger$. Conversely, suppose that $\mathfrak{M}', \langle x, u \rangle \models (\Diamond_1 \psi)^\dagger$. If $\mathfrak{M}', \langle x, u \rangle \models P_\psi$ then there is $v \in W_x, v <_x u$ with $\mathfrak{M}, \langle x, v \rangle \models \psi$. Therefore, there is $v \in W_x, v \neq u$ with $\mathfrak{M}, \langle x, v \rangle \models \psi$. If $\mathfrak{M}', \langle x, u \rangle \models \Diamond_1 \psi^\dagger$ then there is $v \in W_x, v <_x u$ with $\mathfrak{M}', \langle x, v \rangle \models \psi^\dagger$, and so there is $v \in W_x, v \neq u$ with $\mathfrak{M}, \langle x, v \rangle \models \psi$ by the IH. So in both cases $\mathfrak{M}, \langle x, u \rangle \models \Diamond_1 \psi$ follows.

Second, we claim that $\mathfrak{M}', \langle r_0, r_1 \rangle \models \chi_\phi$. Indeed, take any $x \in W$. As r_1 is $<_x$ -least in W_x , we have $\mathfrak{M}', \langle x, r_1 \rangle \models \neg P_\psi$. Now take any $y \in W_x$ with $\mathfrak{M}', \langle x, y \rangle \models \psi^\dagger$ and suppose that $y <_x z$ for some $z \in W_x$. By (73), we have $\mathfrak{M}, \langle x, y \rangle \models \psi$ and so $\mathfrak{M}', \langle x, z \rangle \models P_\psi$ by the definition of \mathfrak{M}' . Finally, suppose that $\mathfrak{M}', \langle x, r_1 \rangle \models \Diamond_1 P_\psi$. Therefore, the set $\{w \in W_x : \langle x, w \rangle \in \mu'(P_\psi)\}$ is non-empty. Let y be its $<_x$ -least element. So there is $z \in W_x, z <_x y$ such that $\mathfrak{M}, \langle x, z \rangle \models \psi$ and $\langle x, z \rangle \notin \mu'(P_\psi)$. Thus $\mathfrak{M}', \langle x, z \rangle \models \neg P_\psi \wedge \psi^\dagger$ by (73). As either $r_1 = y$ or $r_1 <_x y$, we have $\mathfrak{M}', \langle x, r_1 \rangle \models \Diamond_1^+(\neg P_\psi \wedge \psi^\dagger)$, as required.

\Leftarrow : Suppose that $\mathfrak{M}, \langle r_0, r_1 \rangle \models \chi_\phi \wedge \phi^\dagger$ in some model $\mathfrak{M} = \langle \mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}}, \mu \rangle$ based on an expanding 2-frame $\mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}}}$ where $\mathfrak{F} = \langle W, R \rangle$ is transitive and for every $x \in W$, $\mathfrak{G}_x = \langle W_x, <_x \rangle$ is a linear order. Then $W_x \subseteq W_y$ and $<_x \subseteq <_y$ whenever $xRy, x, y \in W$. We may assume that r_0 is a root in \mathfrak{F} , and so $r_1 \in W_x$ for all $x \in W$. Moreover, we may also assume that r_1 is a root in $\langle W_x, <_x \rangle$ for every $x \in W$. Let $\Sigma'_x = \langle W_x, \neq \rangle$, for $x \in W$. Then clearly $\mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}'}} \in \mathcal{C} \times {}^e \mathcal{C}_{diff}$. We define a model $\mathfrak{M}' = \langle \mathfrak{H}_{\mathfrak{F}, \overline{\mathfrak{G}'}} , \mu' \rangle$ by taking $\mu'(P) := \mu(P)$ for all $P \in sub \phi$.

We show by induction on ψ that for all $\psi \in sub \phi, x \in W, u \in W_x$,

$$\mathfrak{M}, \langle x, u \rangle \models \psi^\dagger \quad \text{iff} \quad \mathfrak{M}', \langle x, u \rangle \models \psi. \quad (74)$$

Again, the only interesting case is that of \Diamond_1 . Suppose first that $\mathfrak{M}, \langle x, u \rangle \models (\Diamond_1 \psi)^\dagger$. If

$$\mathfrak{M}, \langle x, u \rangle \models P_\psi, \quad (75)$$

then $r_1 <_x u$ by the first conjunct of χ_ϕ , and so $\mathfrak{M}, \langle x, r_1 \rangle \models \Diamond_1 P_\psi$. So $\mathfrak{M}, \langle x, r_1 \rangle \models \Diamond_1^+(\neg P_\psi \wedge \psi^\dagger)$ follows by the third conjunct of χ_ϕ . So there is $v \in W_x$ with $\mathfrak{M}, \langle x, v \rangle \models \neg P_\psi \wedge \psi^\dagger$, and so $v \neq u$ by (75). Also, by the IH, we have $\mathfrak{M}', \langle x, v \rangle \models \psi$, and so $\mathfrak{M}', \langle x, v \rangle \models \Diamond_1 \psi$ follows as required. The other case when $\mathfrak{M}, \langle x, u \rangle \models \Diamond_1 \psi^\dagger$ is straightforward.

Conversely, suppose that $\mathfrak{M}', \langle x, u \rangle \models \Diamond_1 \psi$. Then there is $v \in W_x, v \neq u$ with $\mathfrak{M}', \langle x, v \rangle \models \psi$, and so by the IH, $\mathfrak{M}, \langle x, v \rangle \models \psi^\dagger$. If $u <_x v$ then $\mathfrak{M}, \langle x, u \rangle \models \Diamond_1 \psi^\dagger$ follows. If $v <_x u$ then by the second conjunct of χ_ϕ , we have $\mathfrak{M}, \langle x, v \rangle \models \Box_1 P_\psi$, and so $\mathfrak{M}, \langle x, u \rangle \models P_\psi$ follows. \square

Next, we show that $\{\langle \omega, < \rangle\} \times {}^e \mathcal{C}_{diff}$ -satisfiability has the ‘finite expanding second components property’:

CLAIM 5.19. *For any formula ϕ , if ϕ is $\{\langle \omega, < \rangle\} \times {}^e \mathcal{C}_{diff}$ -satisfiable, then ϕ is $\{\langle \omega, < \rangle\} \times {}^e \mathcal{C}_{diff}^{fin}$ -satisfiable.*

Proof. Suppose $\mathfrak{M}, \langle 0, r \rangle \models \phi$ for some model \mathfrak{M} based on an expanding 2-frame $\mathfrak{H}_{\langle \omega, < \rangle, \overline{\mathfrak{G}}}$ where $\mathfrak{G}_n = \langle W_n, \neq \rangle$ are difference frames, for $n < \omega$. For all $n < \omega$, $X \subseteq W_n$, we define $cl_n(X)$ as the smallest set Y such that $X \subseteq Y \subseteq W_n$ and having the following property: If $x \in Y$ and $\mathfrak{M}, \langle n, x \rangle \models \Diamond_1 \psi$ for some $\psi \in sub \phi$, then there is $y \in Y$ such that $y \neq x$ and $\mathfrak{M}, \langle n, y \rangle \models \psi$. It is not hard to see that if X is finite then $|cl_n(X)| \leq |X| + 2|sub \phi|$. Now define $\mathfrak{G}'_n := \langle W'_n, \neq \rangle$ by taking $W'_0 := cl_0(\{r\})$ and $W'_{n+1} := cl_{n+1}(W'_n)$ for $n < \omega$. Let \mathfrak{M}' be the restriction of \mathfrak{M} to the expanding 2-frame $\mathfrak{H}_{\langle \omega, < \rangle, \overline{\mathfrak{G}'}}$. A straightforward induction shows that for all $\psi \in sub \phi, n < \omega, w \in W'_n$, we have $\mathfrak{M}, \langle n, w \rangle \models \psi$ iff $\mathfrak{M}', \langle n, w \rangle \models \psi$. \square

Now Theorems 5.14 and 5.16, respectively, follow from Claims 5.18, 5.19 and the following results:

- [27, Thm.1] $\{\langle \omega, < \rangle\} \times {}^e \mathcal{C}_{lin}^{fin}$ -satisfiability is co-r.e.
- [16, Thm.1] $\mathcal{C}_{lin}^{fin} \times {}^e \mathcal{C}_{lin}$ -satisfiability is decidable.

6 Open problems

Our results identify a limit beyond which the one-variable fragment of first-order linear temporal logic is no longer decidable. We have shown that —unlike in the case of the two-variable fragment of classical first-order logic— the addition of even limited counting capabilities ruins decidability in most cases: The resulting logic FOLTL^\neq is very complex over various classes of linear orders, whenever the models have constant, decreasing, or expanding domains. By generalising our techniques to the propositional bimodal setting, we have shown that the bimodal logic $[\mathbf{K4.3}, \mathbf{Diff}]$ of commuting weak order and pseudo-equivalence relations is undecidable. Here are some related unanswered questions:

1. Is the bimodal logic $[\mathbf{K4}, \mathbf{Diff}]$ of commuting transitive and pseudo-equivalence relations decidable? Is the product logic $\mathbf{K4} \times \mathbf{Diff}$ decidable? As $\mathbf{K4}$ can be seen as a notational variant of the fragment of branching time logic CTL that allows only two temporal operators $E\Diamond_F$ and its dual $A\Box_F$, there is another reformulation of the second question: Is the one-variable fragment of first-order CTL decidable when extended with counting and when only $E\Diamond_F$ and $A\Box_F$ are allowed as temporal operators? Note that without counting this coincides with $\mathbf{K4} \times \mathbf{S5} = [\mathbf{K4}, \mathbf{S5}]$ -satisfiability, and that is shown to be decidable by Gabbay and Shehtman [14].
2. Is FOLTL^\neq -satisfiability recursively enumerable in expanding domain models over the class of all linear orders? The bimodal reformulation of this question: Is $\mathcal{C}_{lin} \times {}^e\mathcal{C}_{diff}$ -satisfiability recursively enumerable? By Cor. 5.13, a positive answer would imply decidability of these. Is FOLTL^\neq -satisfiability decidable in expanding domain models over $\langle \mathbb{Q}, < \rangle$ or $\langle \mathbb{R}, < \rangle$?
3. In decreasing 2-frames only ‘half’ of commutativity ($\Box_1\Box_0P \rightarrow \Box_0\Box_1P$) is valid. While in Theorem 4.16 we generalised Theorem 4.1 to classes of decreasing 2-frames and showed that $\mathcal{C}_{lin} \times {}^d\mathcal{C}_{diff}$ -satisfiability is undecidable, it is not clear whether the same can be done in the ‘abstract’ setting: Is satisfiability undecidable in the class of 2-frames having half-commuting weak order and pseudo-equivalence relations?

In our lower bound proofs we used reductions of counter machine problems. Other lower bound results about bimodal logics with grid-like models use reductions of tiling or Turing machine problems [39, 13, 15]. On the one hand, it is not hard to re-prove the same results using counter machine reductions. On the other, it seems tiling and Turing machine techniques require more control over the $\omega \times \omega$ -grid than the limited expressivity that FOLTL^\neq provides. In order to understand the boundary of each technique, it would be interesting to find tiling or Turing machine reductions for the results of this paper.

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